

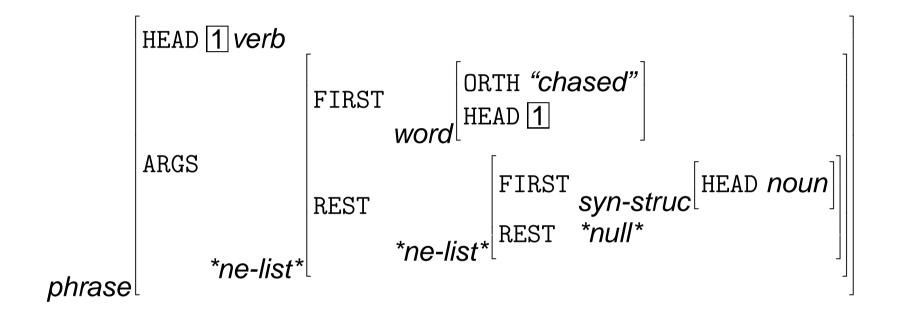
Algorithms for AI and NLP (INF4820 — Unification)

$$\begin{array}{c} \left[\begin{array}{ccc} \text{HEAD} & \boxed{1} \\ \text{SPR} & \langle \rangle \\ \text{COMPS} & \boxed{3} \end{array} \right] & \longrightarrow & \boxed{2} \left[\begin{array}{ccc} \text{SPR} & \langle \rangle \\ \text{COMPS} & \langle \rangle \end{array} \right], & \begin{bmatrix} \text{HEAD} & \boxed{1} \\ \text{SPR} & \left\langle \boxed{2} \right\rangle \\ \text{COMPS} & \boxed{3} \end{array} \right] \\ phrase \end{array}$$

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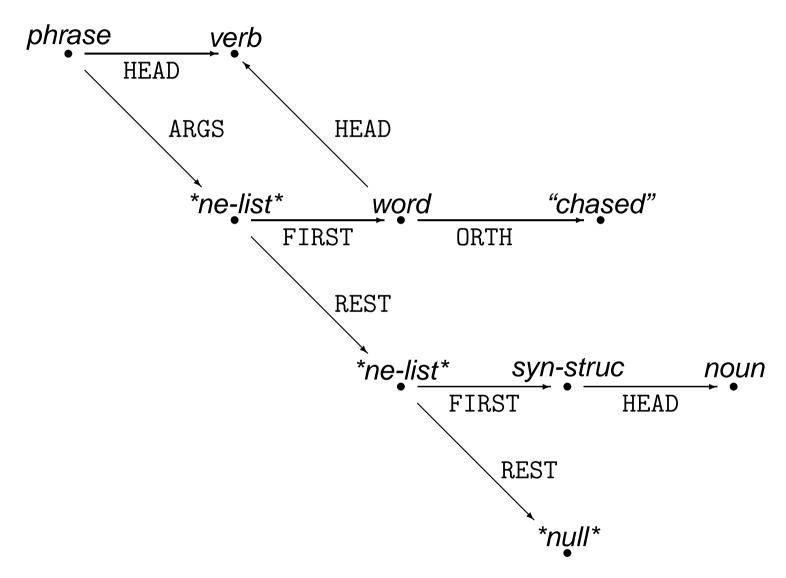
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Feature Structure Reentrancy (AVM)





Feature Structure Reentrancy (DAG)





Feature Structure Unification & Copying

Basic Notions

- Typed feature structures encoded as directed acyclic graphs (DAGs);
- each node bears a type and a set of arcs (aka features value pairs);
- feature structure reentrancy (coreference) corresponds to DAG identity;
- unification creates equivalence classes, encoded through forwarding.

Basic Operations

- *unify()*—make two DAGs equivalent, check and combine all information;
- → at each node, glb() types, forward, recurse over and accumulate arcs;
 - copy()—create structurally equivalent copy (preserving reentrancies);
- → at each node, *copy* slot as short-term memory, reset upon completion.



A Practical Example: Unifying Two DAGs

$$foo^{egin{array}{c} A \ 1 \\ C \ 1 \end{bmatrix}} bar^{egin{array}{c} B \ baz \end{bmatrix}}$$

$$\int_{foo}^{A} bar \begin{bmatrix} B \ baz \end{bmatrix} dt$$



The Costs of Feature Structure Manipulation

Basic Cost Measure

- Visit each DAG node once (node operations 'constant'): full traversal;
- linear in the number of nodes → upper bound is size of largest DAG.

Naïve Complexity Theory

- Prior to each (destructive) unification, make copies of both input DAGs;
- upon completion of each copy, recursively reset copy slot on all nodes.

restore()	copy()	unify()
1	2	5



The unify() vs. copy() Trade-Off

Destructive Unification [Boyer & Moore, 1972]

- Permanently alter both input dags: setf() on forward, type, and arcs;
- → over copying two full copies required for only one result structure;
- → early copying majority of unifications fail: many unnecessary copies.

Non-Destructive Unification [Wroblewski, 1987]

- Incrementally build up result DAG during unification, one node at a time;
- \rightarrow eliminates over copying, reduces early copying more or less effectively.

Quasi-Destructive Unification [Tomabechi, 1991]

- Alter input DAGs in way that is reversible (at small cost): 'generations';
- → copy out result only after unification success, no over or early copying.



Generation Counting

- Protect DAG slots with generation counter → 'expiration date' of value;
- access: require valid generation; assignment: set value and generation;
- → implemented through interaction of global counter and ADT functionality.

```
(defstruct dag
  forward type arcs xcopy (generation 0))
(defparameter *generation* 1)
(defun dag-copy (dag)
   (when (= (dag-generation dag) *generation*) (dag-xcopy dag)))
(defsetf dag-copy dag-set-copy)
(defun dag-set-copy (dag value)
   (setf (dag-generation dag) *generation*)
   (setf (dag-xcopy dag) value))
```



Unification-Based Grammar: Structured Categories

- All (constituent) categories in the grammar are typed feature structures;
- feature structures are recursive, record-like objects: attribute value sets;
- specific TFS configurations may correspond to 'traditional' categories;
- labels like 'S' or 'NP' are mere abbreviations, not elements of the theory.



Interaction of Lexicon and Phrase Structure Schemata

$$\begin{bmatrix} \texttt{HEAD} & \textit{noun} \\ \texttt{SPR} & \langle \, \rangle \\ \texttt{COMPS} & \langle \, \rangle \end{bmatrix}$$

$$\begin{bmatrix} \texttt{HEAD} & \textit{verb} \\ & & \begin{bmatrix} \texttt{HEAD} & \textit{noun} \\ \texttt{SPR} & \left\langle \right\rangle \\ \texttt{COMPS} & \left\langle \right\rangle \end{bmatrix}$$

"snored"



Unification-Based Parsing

Adaptations to CFG-Based Chart Parser

- Make all elements of Σ , C, and P from the grammar feature structures;
- substitute *unification* and *equivalence test* for category comparison;
- unify category of passive edges with argument position of active edges;
- → edge structure LHS is DAG, RHS list of paths to argument positions;
- → fundamental-rule() result of unification is category for new edge;
- → pack-edge() equivalence test: two DAGs contain same information;
 - test *spanning* passive edges for compatibility against start symbol *S*.

```
\#E[id: (i-j) dag --> edge_1 ... edge_i . path_{i+1} ... path_n { alternates }^*]
```

#E[42: (0-8) head-specifier-rule --> 13 . (ARGS REST FIRST)]



Unification-Based Parsing—Practical Concerns

Observations

- Typical systems: 90⁺ per cent of parsing time go to DAG manipulation;
- most unifications fail: predict unification failure cheaply, where possible;
- → rule filter: rule feeding relations; quick check: most likely failure paths;
 - lexicalisation: argument positions in rules may be highly underspecified;
- → head-driven parsing: instantiate RHS bidirectionally, starting from head;
 - many unifications fail very early: copy() more expensive than unify();
- → memory is expensive: redo a couple of unfications instead of one copy.

Several orders of magnitude average speed-up by reducing constants



An Example: The LinGO English Resource Grammar



Suggested Background Activities

- Retrieve the model solution for the fourth exercise from the course site;
- compare our solution to your submission; how is ours better (or not)?
- read [Wroblewski, 1987], be sure to understand over and early copying;
- investigate a call counting scheme for the DAG manipulation routines;
- identify the parts of our earlier parser that require modifications now;
- [Oepen, Flickinger, Tsujii, & Uszkoreit, 2002], Chapters Five and Nine—see whether you can enjoy reading contemporary parsing literature.

