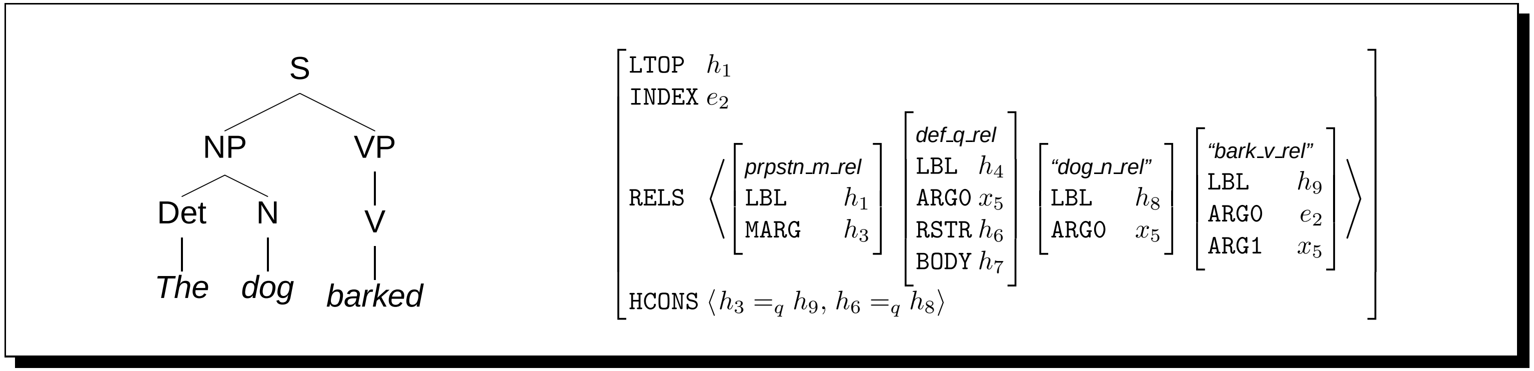

Computational Linguistics (INF2820 — TFSs)

$$\mathit{phrase}\begin{bmatrix}\mathrm{HEAD} & \boxed{1}\\ \mathrm{SPR} & \langle\rangle\\ \mathrm{COMPS} & \boxed{3}\end{bmatrix} \longrightarrow \boxed{2}\,\mathit{phrase}\begin{bmatrix}\mathrm{SPR} & \langle\rangle\\ \mathrm{COMPS} & \langle\rangle\end{bmatrix},\quad \mathit{phrase}\begin{bmatrix}\mathrm{HEAD} & \boxed{1}\\ \mathrm{SPR} & \langle\boxed{2}\rangle\\ \mathrm{COMPS} & \boxed{3}\end{bmatrix}$$

Stephan Oepen

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A Really Complicated Language

[...] *omdat* *ik* *Henk* *de* *nijlpaarden* *zag* *voeren* *.*

A Really Complicated Language

[...] *omdat ik Jan Henk de nijlpaarden zag helpen voeren .*

Limitations of Context-Free Grammar

Agreement and Valency (For Example)

That dog barks.

**That dogs barks.*

**Those dogs barks.*

The dog chased a cat.

**The dog barked a cat.*

**The dog chased.*

**The dog chased a cat my neighbours.*

The cat was chased by a dog.

**The cat was chased of a dog.*

...

Agreement and Valency in Context-Free Grammars

Grammatical Functions

Licensing — Government — Agreement

*The dog barks. — *The dog a cat barks — *The dog barks a cat.*

*Kim depends on Sandy — *Kim depends in Sandy*

The class meets on Thursday in 3A at 14:15.

- **Constituent** node in analysis tree (lexical entry or instance of rule);
- **Head** licenses additional constituents and can govern their form;
- **Specifier** precedes head, singleton, nominative case, agreement;
- **Complement** post-head, licensed and governed, order constraints;
- **Adjunct** 'free' modifier, optional, may iterate, designated position;
- **Government** directed: a property of c_1 determines the form of c_2;
- **Agreement** bi-directional: co-occurence of properties on c_1 and c_2.

Structured Categories in a Unification Grammar

- All (constituent) categories in the grammar are typed feature structures;
- specific TFS configurations may correspond to ‘traditional’ categories;
- labels like ‘S’ or ‘NP’ are mere abbreviations, not elements of the theory.

$$\mathit{word}\begin{bmatrix} \text{HEAD} & \mathit{noun} \\ \text{SPR} & \left\langle \begin{bmatrix} \text{HEAD}\ \mathit{det} \end{bmatrix} \right\rangle \\ \text{COMPS} & \langle\rangle \end{bmatrix} \qquad \mathit{phrase}\begin{bmatrix} \text{HEAD} & \mathit{verb} \\ \text{SPR} & \langle\rangle \\ \text{COMPS} & \langle\rangle \end{bmatrix} \qquad \mathit{phrase}\begin{bmatrix} \text{HEAD} & \mathit{verb} \\ \text{SPR} & \left\langle \begin{bmatrix} \text{HEAD}\ \mathit{noun} \end{bmatrix} \right\rangle \\ \text{COMPS} & \langle\rangle \end{bmatrix}$$

‘N’	‘S’	‘VP’
‘lexical’	‘maximal’	‘intermediate’

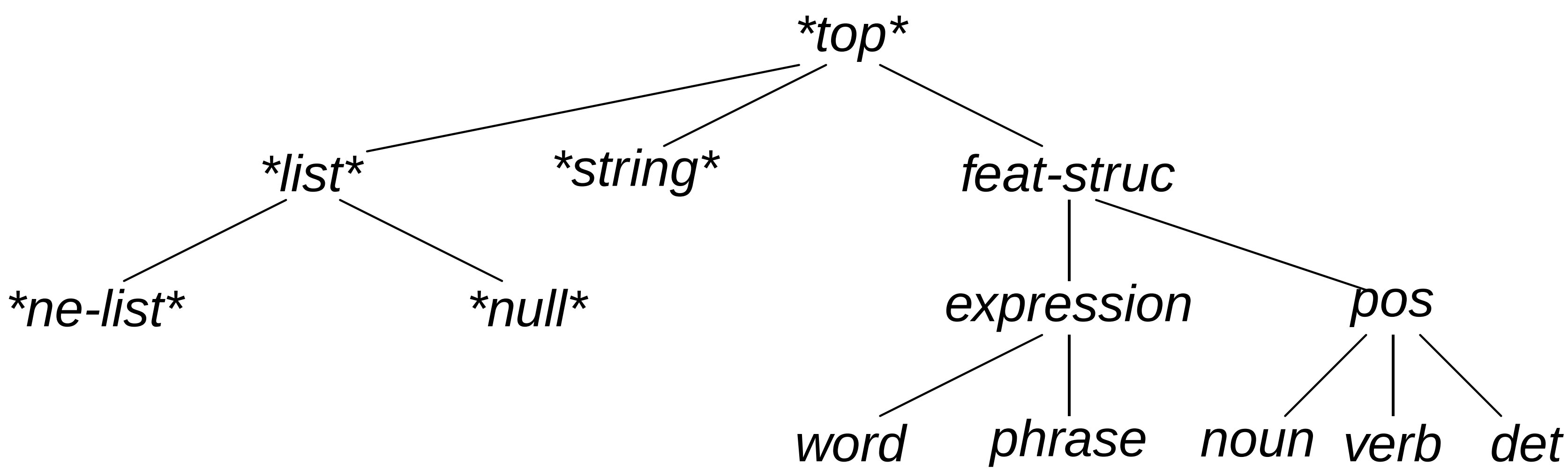
Interaction of Lexicon and Phrase Structure Schemata

$$\textit{phrase}\begin{bmatrix}\text{HEAD} & \boxed{1}\\ \text{SPR} & \langle\,\rangle\\ \text{COMPS} & \boxed{3}\end{bmatrix} \longrightarrow \boxed{2}\,\textit{phrase}\begin{bmatrix}\text{SPR} & \langle\,\rangle\\ \text{COMPS} & \langle\,\rangle\end{bmatrix},\quad \textit{phrase}\begin{bmatrix}\text{HEAD} & \boxed{1}\\ \text{SPR} & \langle\boxed{2}\rangle\\ \text{COMPS} & \boxed{3}\end{bmatrix}$$

$$\begin{bmatrix}\text{ORTH} & \textit{“Kim”}\\ \text{HEAD} & \textit{noun}\\ \text{SPR} & \langle\,\rangle\\ \text{COMPS} & \langle\,\rangle\end{bmatrix} \qquad \begin{bmatrix}\text{ORTH} & \textit{“sleeps”}\\ \text{HEAD} & \textit{verb}\\ \text{SPR} & \left\langle\begin{bmatrix}\text{HEAD} & \textit{noun}\\ \text{SPR} & \langle\,\rangle\\ \text{COMPS} & \langle\,\rangle\end{bmatrix}\right\rangle\\ \text{COMPS} & \langle\,\rangle\end{bmatrix}$$

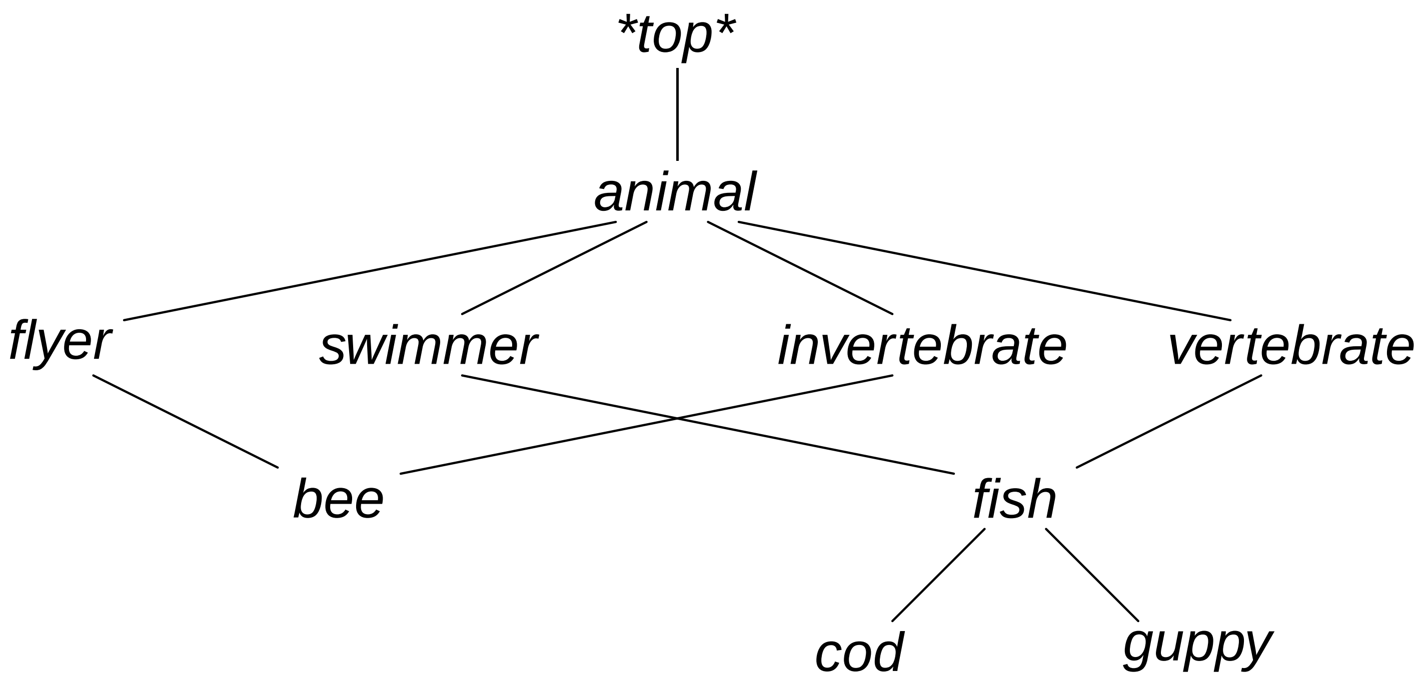
The Type Hierarchy: Fundamentals

- Types ‘represent’ groups of entities with similar properties (‘classes’);
- types ordered by specificity: subtypes inherit properties of (all) parents;
- type hierarchy determines which types are compatible (and which not).



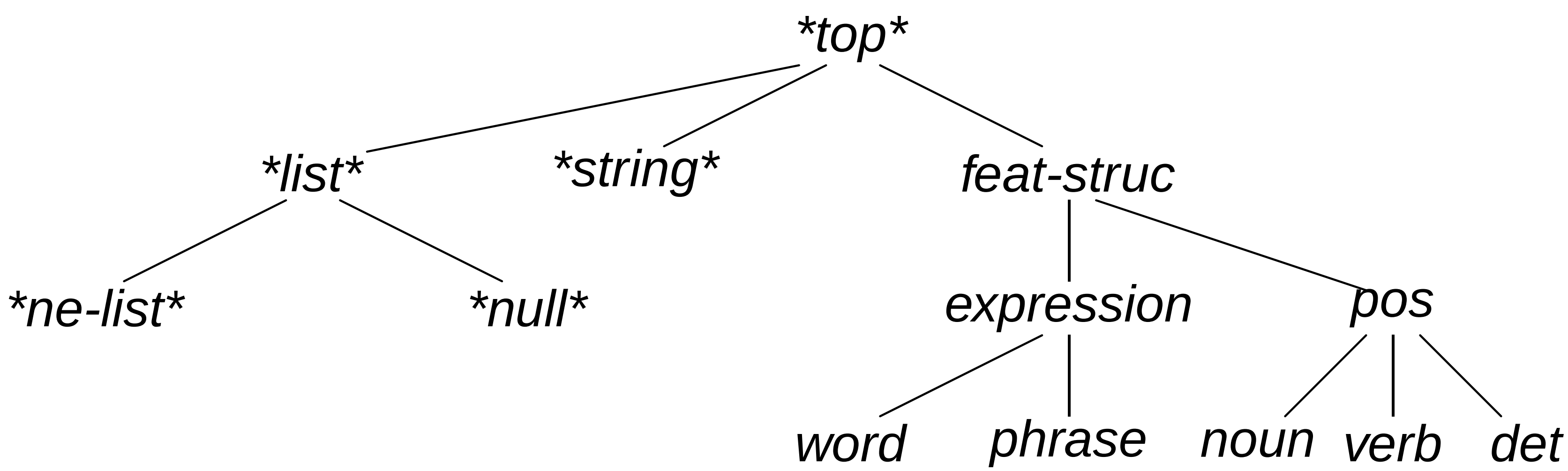
Multiple Inheritance

- *flyer* and *swimmer* no common descendants: they are incompatible;
- *flyer* and *bee* stand in hierarchical relationship: they unify to subtype;
- *flyer* and *invertebrate* have a unique greatest common descendant.



The Type Hierarchy — Appropriate Features

- Features record properties of entities; in turn, feature values are TFSs;
- features are defined by a unique most general type: appropriateness;
- feature values constrained to a specific type → monotonic inheritance.



Typed Feature Structure Subsumption

- Typed feature structures can be partially ordered by information content;
- a more general structure is said to *subsume* a more specific one;
- $_{*top*}\Big[\;\Big]$ is the most general feature structure (while $\bot$ is inconsistent);
- $\sqsubseteq$ ('square subset or equal') conventionally used to depict subsumption.

Feature structure F subsumes feature structure G ($F \sqsubseteq G$) iff: (1) if path p is defined in F then p is also defined in G and the type of the value of p in F is a supertype or equal to the type of the value of p in G, and (2) all paths that are reentrant in F are also reentrant in G.

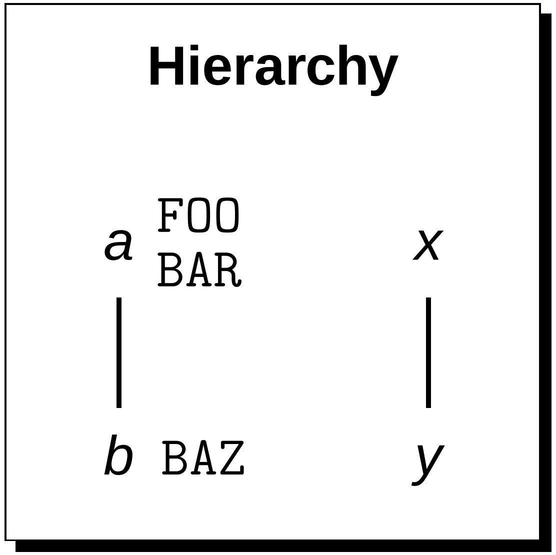
Feature Structure Subsumption: Examples

TFS$_1$: $\begin{bmatrix} a \\ \text{FOO } x \\ \text{BAR } x \end{bmatrix}$

TFS$_2$: $\begin{bmatrix} a \\ \text{FOO } x \\ \text{BAR } y \end{bmatrix}$

TFS$_3$: $\begin{bmatrix} b \\ \text{FOO } y \\ \text{BAR } x \\ \text{BAZ } x \end{bmatrix}$

TFS$_4$: $\begin{bmatrix} a \\ \text{FOO } \boxed{1} x \\ \text{BAR } \boxed{1} \end{bmatrix}$



Feature structure F subsumes feature structure G ($F \sqsubseteq G$) iff: (1) if path p is defined in F then p is also defined in G and the type of the value of p in F is a supertype or equal to the type of the value of p in G, and (2) all paths that are reentrant in F are also reentrant in G.

Typed Feature Structure Unification

- Decide whether two typed feature structures are mutually compatible;
- determine combination of two TFSs to give the most general feature structure which retains all information which they individually contain;
- if there is no such feature structure, unification fails (depicted as $\perp$);
- unification *monotonically* combines information from both 'input' TFSs;
- *relation to subsumption* the unification of two structures F and G is the most general TFS which is subsumed by both F and G (if it exists).
- $\sqcap$ ('square set intersection') conventionally used to depict unification.

Typed Feature Structure Unification: Examples

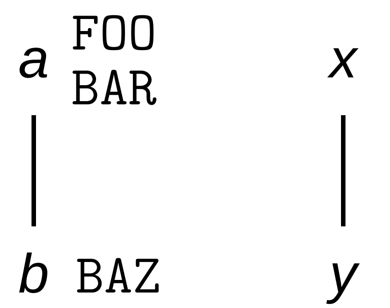
TFS$_1$: $\left[\begin{array}{l} \text{FOO } x \\ \text{BAR } x \end{array}\right]_a$

TFS$_2$: $\left[\begin{array}{l} \text{FOO } x \\ \text{BAR } y \end{array}\right]_a$

TFS$_3$: $\left[\begin{array}{l} \text{FOO } y \\ \text{BAR } x \\ \text{BAZ } x \end{array}\right]_b$

TFS$_4$: $\left[\begin{array}{l} \text{FOO } \boxed{1}\, x \\ \text{BAR } \boxed{1} \end{array}\right]_a$

Hierarchy



$$\text{TFS}_1 \sqcap \text{TFS}_2 \equiv \text{TFS}_2 \qquad \text{TFS}_1 \sqcap \text{TFS}_3 \equiv \text{TFS}_3 \qquad \text{TFS}_3 \sqcap \text{TFS}_4 \equiv \left[\begin{array}{l} \text{FOO } \boxed{1}\, y \\ \text{BAR } \boxed{1} \\ \text{BAZ } x \end{array}\right]_b$$

Type Constraints and Appropriate Features

- Well-formed TFSs satisfy all *type constraints* from the type hierarchy;
- type constraints are typed feature structures associated with a type;
- the top-level features of a type constraint are *appropriate features*;
- type constraints express generalizations over a ‘class’ (set) of objects.

type	constraint	appropriate features
ne-list	**ne-list** [FIRST **top**, REST **list**]	FIRST and REST

Recursion in the Type Hierarchy

- Type hierarchy must be finite *after* type inference; illegal type constraint:

```
*list* := *top* & [ FIRST *top*, REST *list* ].
```

- needs additional provision for empty lists; indirect recursion:

```
*list* := *top*.
*ne-list* := *list* & [ FIRST *top*, REST *list* ].
*null* := *list*.
```

- recursive types allow for *parameterized list types* ('list of X'):

```
*s-list* := *list*.
*s-ne-list* := *ne-list* & *s-list &
               [ FIRST expression, REST *s-list* ].
*s-null* := *null* & *s-list*.
```

Notational Conventions

- lists not available as built-in data type; abbreviatory notation in TDL:

```
< a, b > ≡ [ FIRST a, REST [ FIRST b, REST *null* ] ]
```

- underspecified (variable-length) list:

```
< a, ... > ≡ [ FIRST a, REST *list* ]
```

- difference (open-ended) lists; allow concatenation by unification:

```
<! a !> ≡ [ LIST [ FIRST a, REST #tail ], LAST #tail ]
```

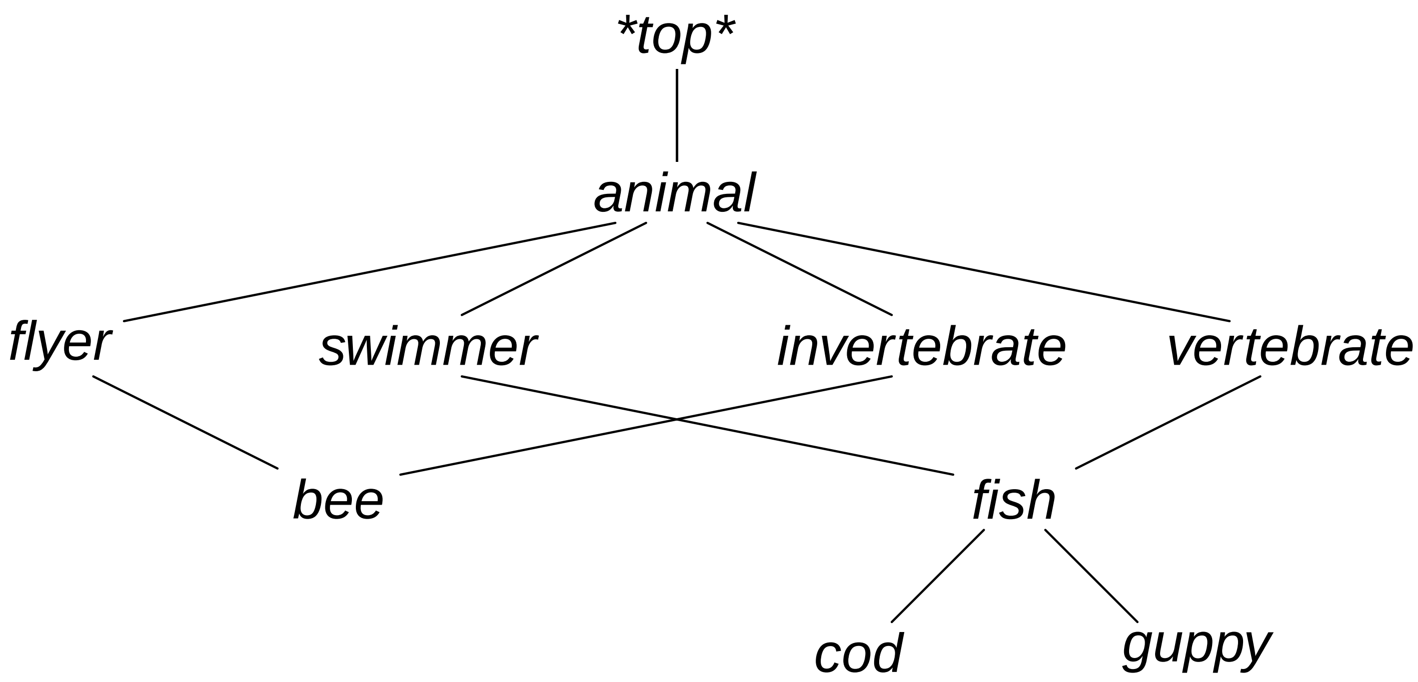
- built-in and 'non-linguistic' types pre- and suffixed by asterisk (**top**);
- strings (e.g. *"chased"*) need no declaration; always subtypes of **string**;
- strings cannot have subtypes and are (thus) mutually incompatible.

Properties of (Our) Type Hierarchies

- **Unique Top** a single hierarchy of all types with a unique top node;
- **No Cycles** no path through the hierarchy from one type to itself;
- **Unique Greatest Lower Bounds** Any two types in the hierarchy are either (a) incompatible (i.e. share no descendants) or (b) have a unique most general ('highest') descendant (called their greatest lower bound);
- **Closed World** all types that exist have a known position in hierarchy;
- **Compatibility** type compatibility in the hierarchy determines feature structure unifiability: two types unify to their greatest lower bound.

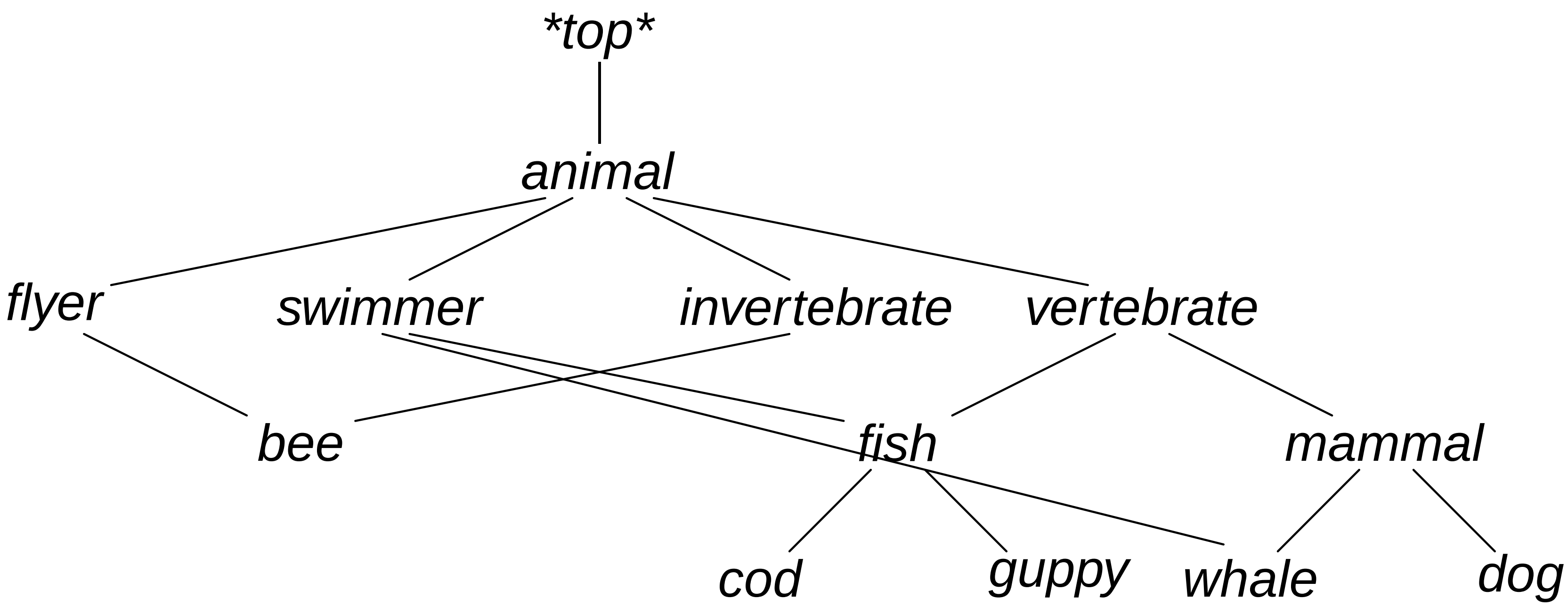
Multiple Inheritance (Repeated for Convenience)

- *flyer* and *swimmer* no common descendants: they are incompatible;
- *flyer* and *bee* stand in hierarchical relationship: they unify to subtype;
- *flyer* and *invertebrate* have a unique greatest common descendant.



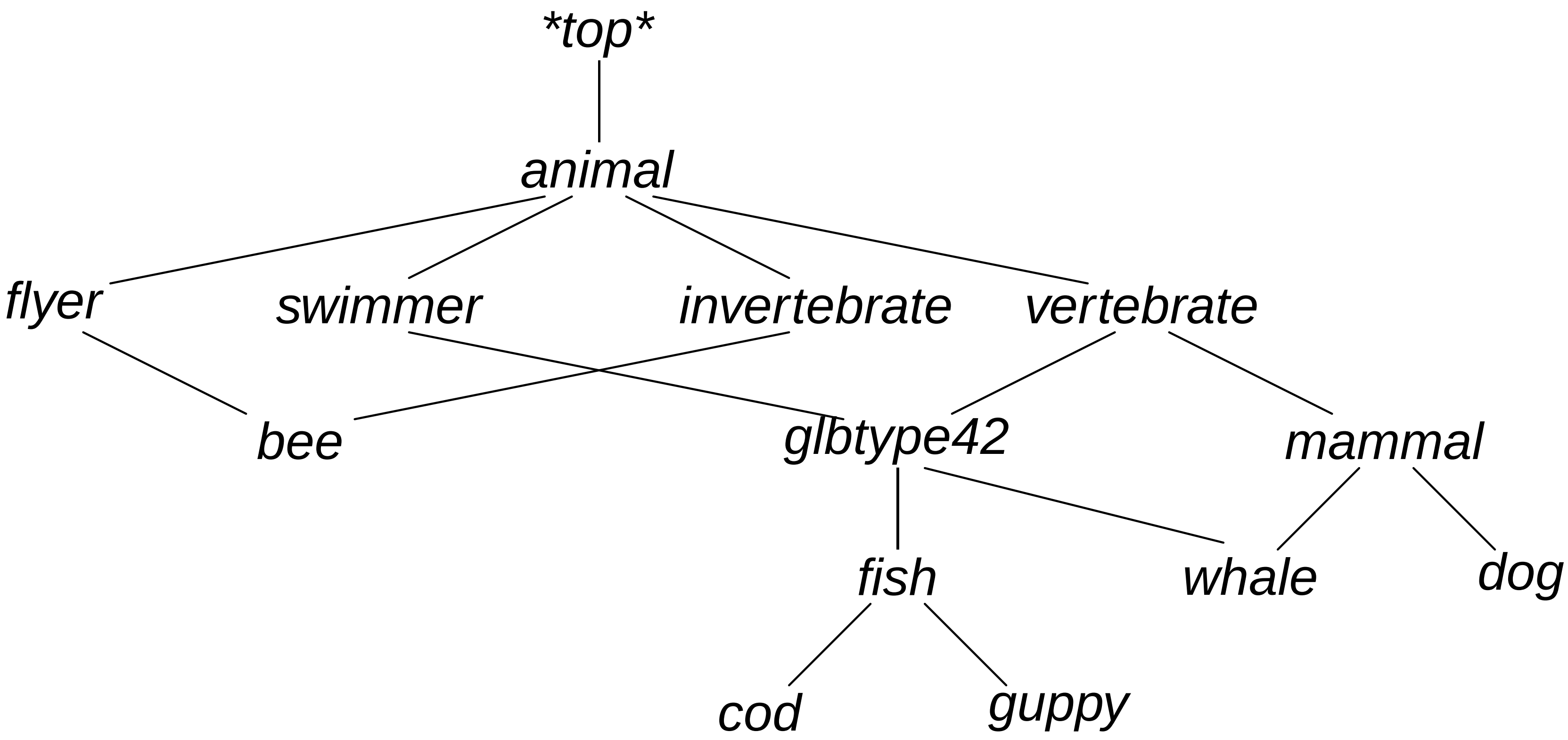
An Invalid Type Hierarchy

- *swimmer* and *vertebrate* have two joint descendants: *fish* and *whale*;
- *fish* and *whale* are incomparable in the hierarchy: glb condition violated.



Fixing the Type Hierarchy

- LKB system introduces *glb types* as required: '*swimmer-vertebrate*'.



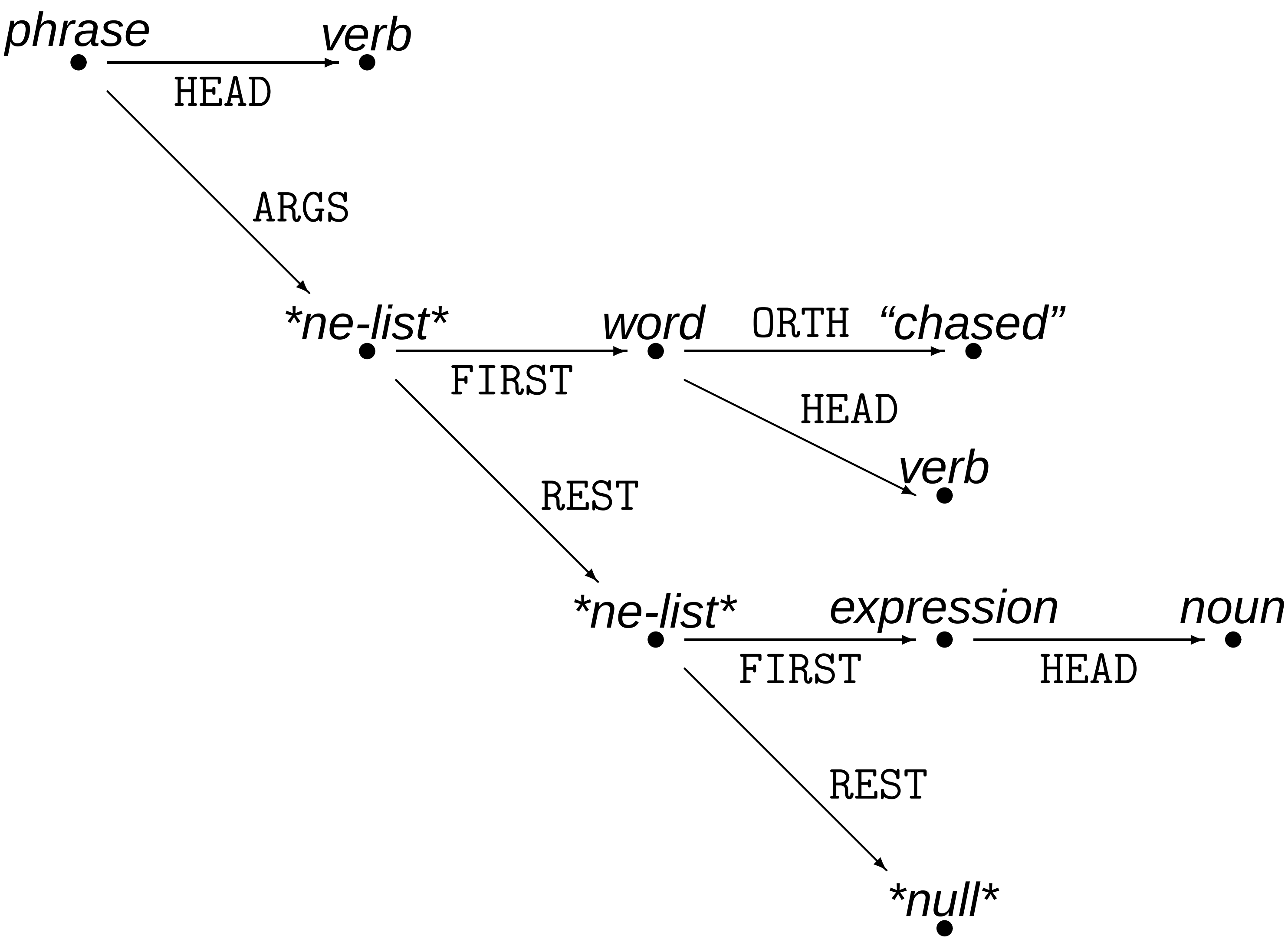
Properties of (Our) Typed Feature Structures

- **Finiteness** a typed feature structure has a finite number of nodes;
- **Unique Root and Connectedness** a typed feature structure has a unique root node; apart from the root, all nodes have at least one parent;
- **No Cycles** no node has an arc that points back to the root node or to another node that intervenes between the node itself and the root;
- **Unique Features** any node can have any (finite) number of outgoing arcs, but the arc labels (i.e. features) must be unique within each node;
- **Typing** each node has single type which is defined in the hierarchy.

Typed Feature Structure Example (as AVM)

```
phrase [ HEAD  verb
         ARGS  *ne-list* [ FIRST  word [ ORTH "chased"
                                         HEAD verb ]
                           REST   *ne-list* [ FIRST  expression [ HEAD noun ]
                                              REST   *null* ] ] ]
```

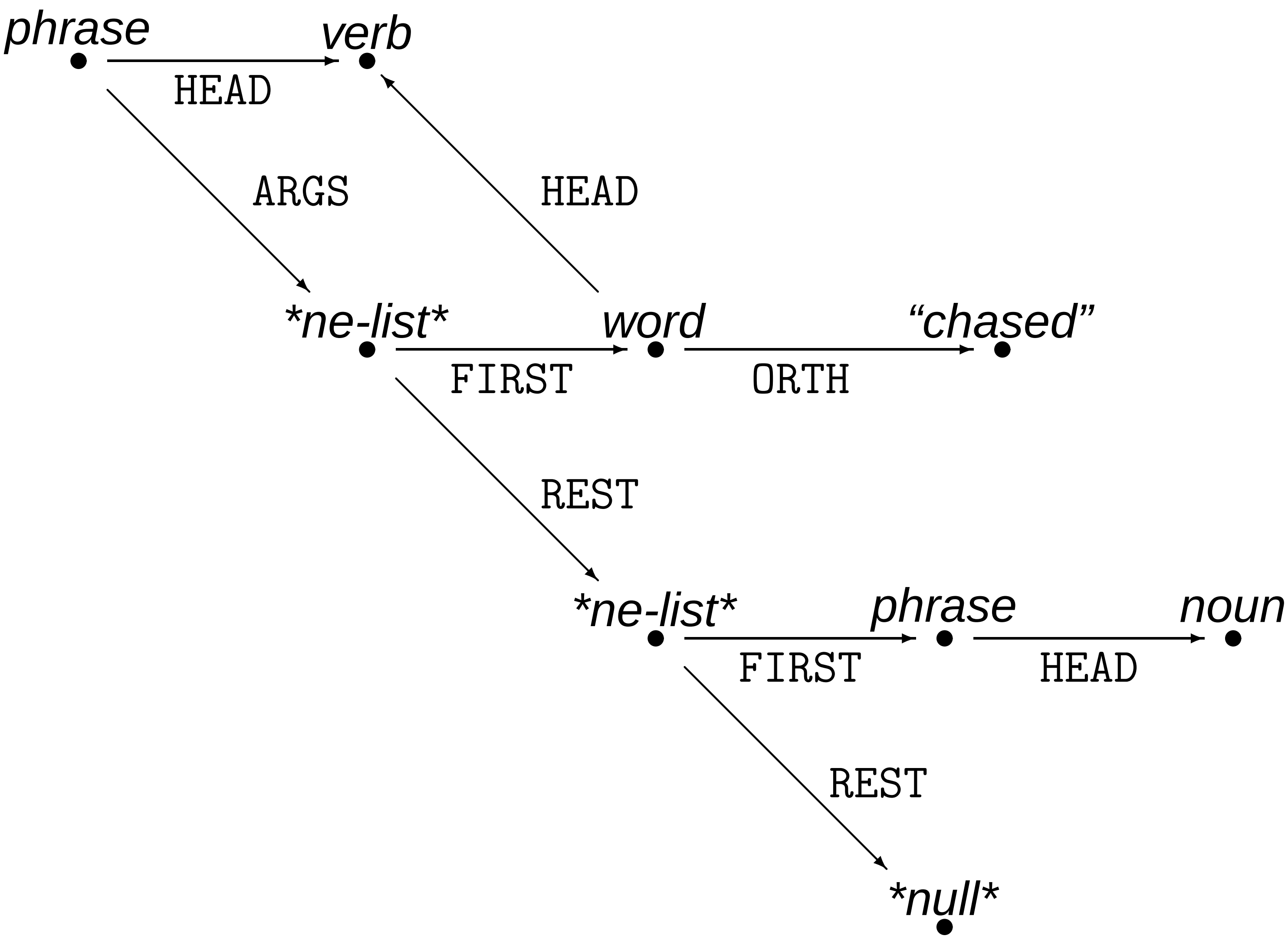
Typed Feature Structure Example (as Graph)



Typed Feature Structure Example (in TDL)

```
vp := phrase &
[ HEAD verb,
  ARGS *ne-list* &
        [ FIRST word &
                 [ ORTH "chased",
                   HEAD verb ],
          REST *ne-list* &
                [ FIRST expression &
                         [ HEAD noun ],
                  REST *null* ]]] .
```

Reentrancy in a Typed Feature Structure (Graph)



Reentrancy in a Typed Feature Structure (AVM)

$$
\mathit{phrase}\begin{bmatrix}
\text{HEAD} & \boxed{1}\,\mathit{verb} \\
\text{ARGS} & \mathit{*ne\text{-}list*}\begin{bmatrix}
\text{FIRST} & \mathit{word}\begin{bmatrix}
\text{ORTH} & \text{“\textit{chased}”} \\
\text{HEAD} & \boxed{1}
\end{bmatrix} \\
\text{REST} & \mathit{*ne\text{-}list*}\begin{bmatrix}
\text{FIRST} & \mathit{phrase}\begin{bmatrix}\text{HEAD} & \mathit{noun}\end{bmatrix} \\
\text{REST} & \mathit{*null*}
\end{bmatrix}
\end{bmatrix}
\end{bmatrix}
$$

Reentrancy in a Typed Feature Structure (TDL)

```
vp := phrase &
[ HEAD #head & verb,
  ARGS *ne-list* &
       [ FIRST word &
               [ ORTH "chased",
                 HEAD #head ],
         REST *ne-list* &
              [ FIRST phrase &
                      [ HEAD noun ],
                REST *null* ]]] .
```

The Linguistic Knowledge Builder (LKB)

Compiler and Interactive Debugger

- Grammar definition errors identified at load time by position in file;
- inheritance and appropriateness tracked by type and attributes;
- batch check, expansion, and indexing of full lexicon on demand;
- efficient parser and generator to map between strings and meaning;
- visualization of main data types; interactive stepping and unification.

- Main developers: Copestake (original), Carroll, Malouf, and Oepen;
- implementation: Allegro CL, Macintosh CL, (LispWorks, CMU CL);
- available in open-source and binary form for common platforms.