

Computational Linguistics (INF2820 — TFSs)

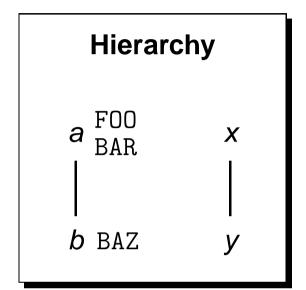
$$\begin{array}{c} \left[\begin{array}{ccc} \text{HEAD} & \boxed{1} \\ \text{SPR} & \langle \rangle \\ \text{COMPS} & \boxed{3} \end{array} \right] & \longrightarrow & \boxed{2} \left[\begin{array}{ccc} \text{SPR} & \langle \rangle \\ \text{COMPS} & \langle \rangle \end{array} \right], & \begin{bmatrix} \text{HEAD} & \boxed{1} \\ \text{SPR} & \left\langle \boxed{2} \right\rangle \\ \text{COMPS} & \boxed{3} \end{array} \right] \\ phrase \end{array}$$

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Feature Structure Subsumption: Examples

TFS₁:
$$\begin{bmatrix} F00 \ X \\ BAR \ X \end{bmatrix}$$
TFS₂: $\begin{bmatrix} F00 \ X \\ BAR \ Y \end{bmatrix}$
TFS₃: $\begin{bmatrix} F00 \ Y \\ BAR \ X \\ BAZ \ X \end{bmatrix}$
TFS₄: $\begin{bmatrix} F00 \ 1 \ X \\ BAR \ 1 \end{bmatrix}$



Feature structure F subsumes feature structure G ($F \sqsubseteq G$) iff: (1) if path p is defined in F then p is also defined in G and the type of the value of p in F is a supertype or equal to the type of the value of p in G, and (2) all paths that are reentrant in F are also reentrant in G.

Typed Feature Structure Unification

- Decide whether two typed feature structures are mutually compatible;
- determine combination of two TFSs to give the most general feature structure which retains all information which they individually contain;
- \bullet if there is no such feature structure, unification fails (depicted as \bot);
- unification *monotonically* combines information from both 'input' TFSs;
- relation to subsumption the unification of two structures F and G is the most general TFS which is subsumed by both F and G (if it exists).
- □ ('square set intersection') conventionally used to depict unification.



Typed Feature Structure Unification: Examples

TFS₁:
$$\begin{bmatrix} F00 & x \\ BAR & x \end{bmatrix}$$
TFS₂:
$$\begin{bmatrix} F00 & x \\ BAR & y \end{bmatrix}$$
TFS₃:
$$\begin{bmatrix} F00 & y \\ BAR & x \\ BAZ & x \end{bmatrix}$$
TFS₄:
$$\begin{bmatrix} F00 & 1 \\ BAR & 1 \end{bmatrix}$$

$$\mathsf{TFS}_1 \sqcap \mathsf{TFS}_2 \equiv \mathsf{TFS}_2 \quad \mathsf{TFS}_1 \sqcap \mathsf{TFS}_3 \equiv \mathsf{TFS}_3 \quad \mathsf{TFS}_3 \sqcap \mathsf{TFS}_4 \equiv \begin{bmatrix} \mathsf{F00} \ \boxed{1} \textit{y} \\ \mathsf{BAR} \ \boxed{1} \\ \mathsf{BAZ} \ \textit{x} \end{bmatrix}$$



Type Constraints and Appropriate Features

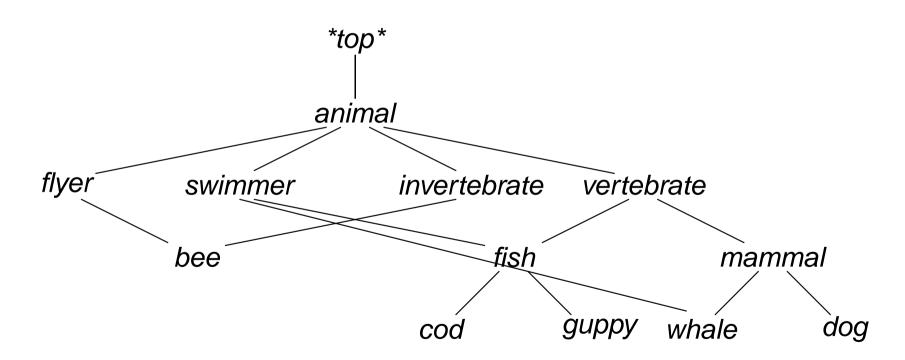
- Well-formed TFSs satisfy all type constraints from the type hierarchy;
- type constraints are typed feature structures associated with a type;
- the top-level features of a type constraint are appropriate features;
- type constraints express generalizations over a 'class' (set) of objects.

type	constraint	appropriate features
ne-list	*ne-list* FIRST *top*	FIRST and REST



An Invalid Type Hierarchy

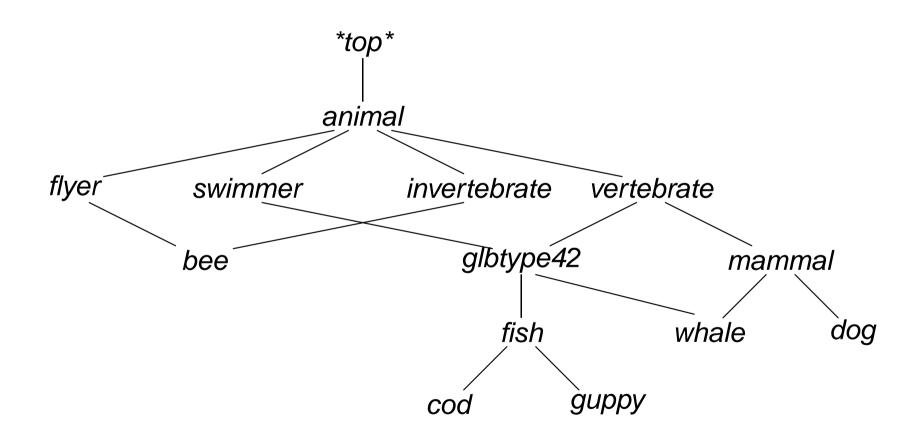
- swimmer and vertebrate have two joint descendants: fish and whale;
- fish and whale are incomparable in the hierarchy: glb condition violated.





Fixing the Type Hierarchy

• LKB system introduces *glb types* as required: 'swimmer-vertebrate'.





More Interesting Well-Formed Unification

Type Constraints Associated to animal Hierarchy

$$swimmer \rightarrow swimmer \begin{bmatrix} \texttt{FINS bool} \end{bmatrix} \qquad mammal \rightarrow mammal \begin{bmatrix} \texttt{FRIENDLY bool} \end{bmatrix}$$

$$whale \rightarrow \begin{bmatrix} \texttt{BALEEN bool} \\ \texttt{FINS true} \\ \texttt{FRIENDLY bool} \end{bmatrix}$$

$$mammal$$
 $\begin{bmatrix} ext{FRIENDLY true} \end{bmatrix} \sqcap \begin{bmatrix} ext{swimmer} \end{bmatrix} \equiv \begin{bmatrix} ext{BALEEN} & bool \\ ext{FINS} & true \\ ext{whale} \end{bmatrix}$ $mammal$ $\begin{bmatrix} ext{FRIENDLY true} \end{bmatrix} \sqcap \begin{bmatrix} ext{Swimmer} \end{bmatrix} \equiv 1$



Recursion in the Type Hierarchy

• Type hierarchy must be finite *after* type inference; illegal type constraint:

```
*list* := *top* & [ FIRST *top*, REST *list* ].
```

• needs additional provision for empty lists; indirect recursion:

```
*list* := *top*.

*ne-list* := *list* & [ FIRST *top*, REST *list* ].

*null* := *list*.
```

• recursive types allow for *parameterized list types* ('list of X'):



Properties of (Our) Type Hierarchies

- Unique Top a single hierarchy of all types with a unique top node;
- No Cycles no path through the hierarchy from one type to itself;
- Unique Greatest Lower Bounds Any two types in the hierarchy are either (a) incompatible (i.e. share no descendants) or (b) have a unique most general ('highest') descendant (called their greatest lower bound);
- Closed World all types that exist have a known position in hierarchy;
- **Compatibility** type compatibility in the hierarchy determines feature structure unifiability: two types unify to their greatest lower bound.



Properties of (Our) Typed Feature Structures

- Finiteness a typed feature structure has a finite number of nodes;
- Unique Root and Connectedness a typed feature structure has a unique root node; apart from the root, all nodes have at least one parent;
- No Cycles no node has an arc that points back to the root node or to another node that intervenes between the node itself and the root;
- Unique Features any node can have any (finite) number of outgoing arcs, but the arc labels (i.e. features) must be unique within each node;
- **Typing** each node a has single type which is defined in the hierarchy.



The Linguistic Knowledge Builder (LKB)

Compiler and Interactive Debugger

- Grammar definition errors identified at load time by position in file;
- inheritance and appropriateness tracked by type and attributes;
- batch check, expansion, and indexing of full lexicon on demand;
- efficient parser and generator to map between strings and meaning;
- visualization of main data types; interactive stepping and unification.

- Main developers: Copestake (original), Carroll, Malouf, and Oepen;
- implementation: Allegro CL, Macintosh CL, (LispWorks, CMU CL);
- available in open-source and binary form for common platforms.



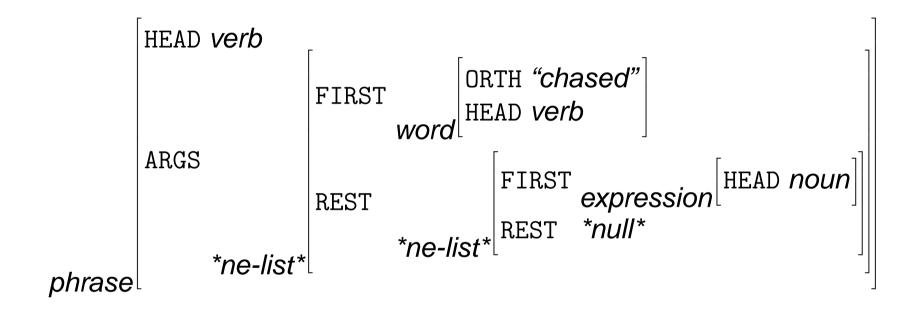
The Format of Grammar Rules in the LKB

$$\begin{array}{c} \left[\begin{array}{c} \text{HEAD} & \boxed{1} \\ \text{SPR} & \langle \rangle \\ \text{COMPS} & \boxed{3} \end{array} \right] & \longrightarrow & \boxed{2} \left[\begin{array}{c} \text{SPR} & \langle \rangle \\ \text{COMPS} & \langle \rangle \end{array} \right], \\ \textit{phrase} \end{array} \\ \begin{array}{c} \text{Phrase} \\ \text{Phrase} \end{array} \right]$$

$$\begin{array}{c|c} \text{HEAD} & \boxed{1} \\ \text{SPR} & \langle \rangle \\ \text{COMPS} & \boxed{3} \\ \\ \text{ARGS} & \langle \boxed{2} \begin{bmatrix} \text{SPR} & \langle \rangle \\ \text{COMPS} & \langle \rangle \end{bmatrix}, \\ \text{phrase} & \begin{array}{c} \text{HEAD} & \boxed{1} \\ \text{SPR} & \langle \boxed{2} \rangle \\ \text{COMPS} & \boxed{3} \end{array} \end{array} \right]$$

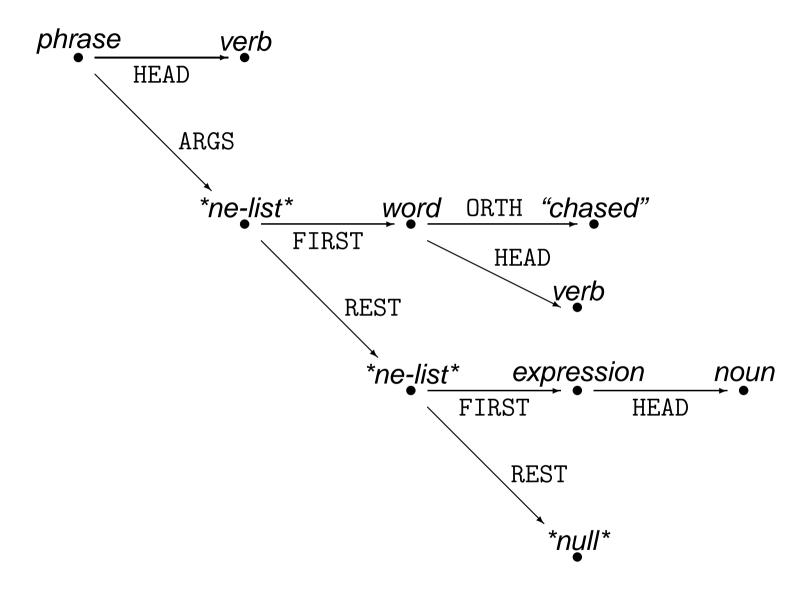


Typed Feature Structure Example (as AVM)





Typed Feature Structure Example (as Graph)

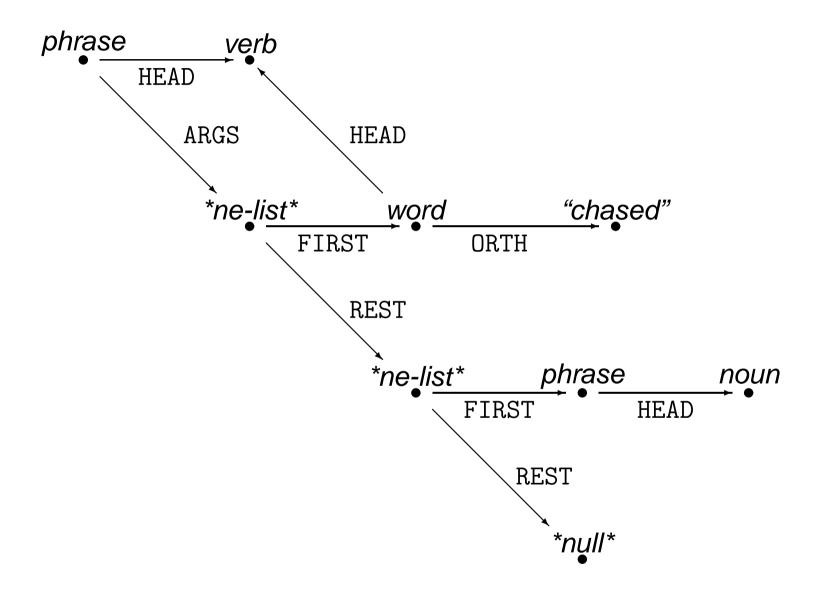




Typed Feature Structure Example (in TDL)

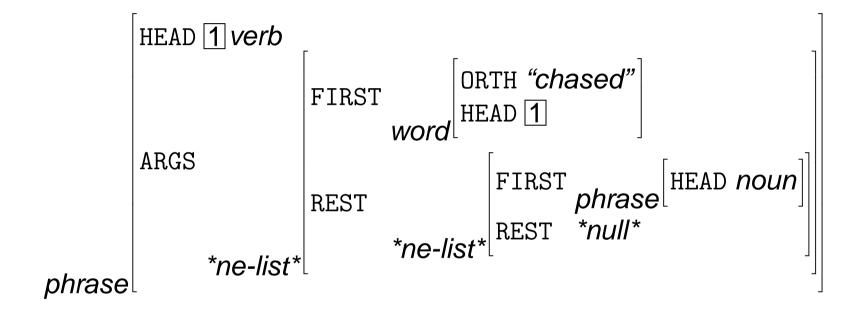


Reentrancy in a Typed Feature Structure (Graph)





Reentrancy in a Typed Feature Structure (AVM)





Reentrancy in a Typed Feature Structure (TDL)



More Terminology: Grammatical Functions

Licensing — Government — Agreement

The dog barks. — *The dog a cat barks — *The dog barks a cat.

Kim depends on Sandy — *Kim depends in Sandy

The class meets on Thursday in 508 at 12:15.

- Constituent node in analysis tree (terminal or instantiation of rule);
- **Head** licenses additional constituents and can govern their form;
- Specifier precedes head, singleton, nominative case, agreement;
- Complement post-head, licensed and governed, order constraints;
- Adjunct 'free' modifier, optional, may iterate, designated position;
- **Government** directed: a property of c_1 determines the form of c_2 ;
- **Agreement** bi-directional: co-occurrence of properties on c_1 and c_2 .



An Ambiguous Example

Kim shoveled snow on lifts.



A Highly Ambiguous Example

The manager placed his bid on my desk.

