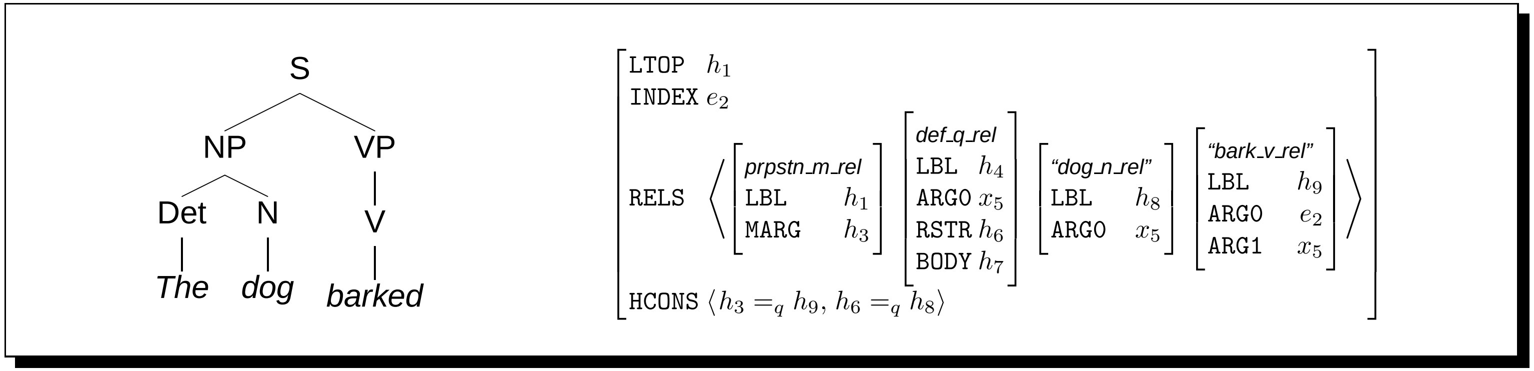

Computational Linguistics (INF2820 — Parsing)

$S \longrightarrow NP\ VP;\ \ S \longrightarrow S\ PP;\ \ S \longrightarrow VP$

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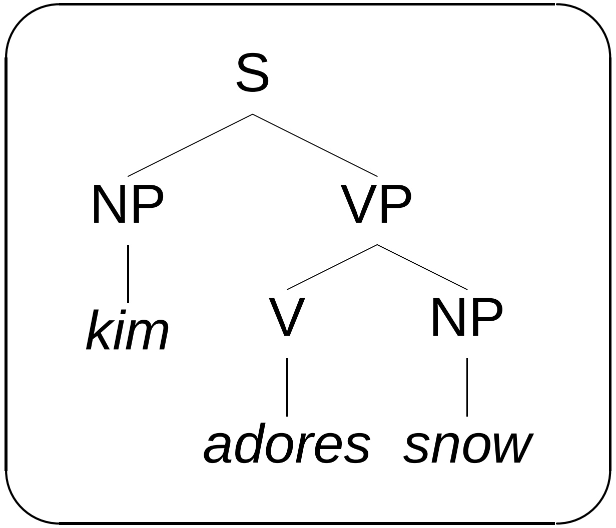
Background: Trees as Bracketed Sequences

- Trees can be encoded as sequences (*dominance* plus *precedence*):

```
(S (NP kim)
   (VP (V adored)
       (NP snow)))
```



- the `first()` element (at each level) represents the tree root (or mother);
- all other elements (i.e. the `rest()`) correspond to immediate daughters.

Mildly Mathematically: Context-Free Grammars

- Formally, a *context-free grammar* (CFG) is a quadruple: $\langle C, \Sigma, P, S \rangle$
- C is the set of categories (aka *non-terminals*), e.g. $\{$S, NP, VP, V$\}$;
- Σ is the vocabulary (aka *terminals*), e.g. $\{$Kim, snow, saw, in$\}$;
- P is a set of category rewrite rules (aka *productions*), e.g.

S $\rightarrow$ NP VP
VP $\rightarrow$ V NP
NP $\rightarrow$ Kim
NP $\rightarrow$ snow
V $\rightarrow$ saw

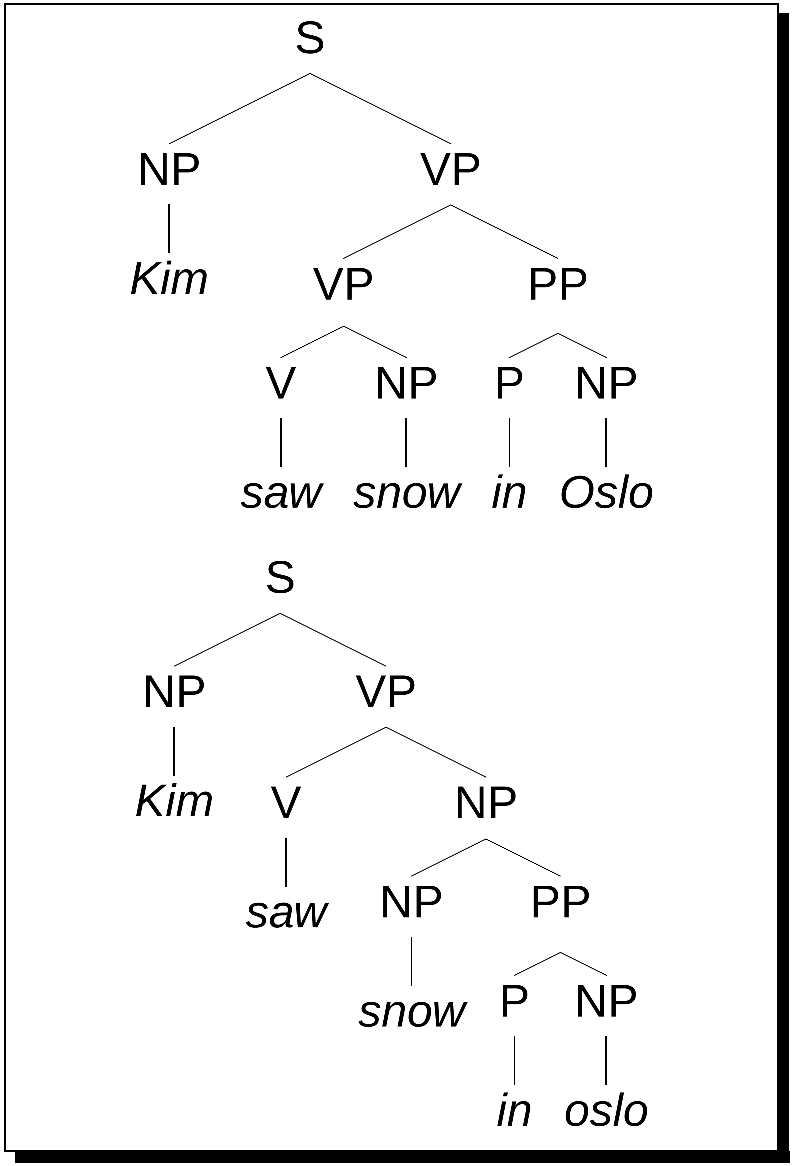
- $S \in C$ is the *start symbol*, a filter on complete ('sentential') results;
- for each rule '$\alpha \rightarrow \beta_1, \beta_2, ..., \beta_n$' $\in P$: $\alpha \in C$ and $\beta_i \in C \cup \Sigma$; $1 \leq i \leq n$.

Parsing: Recognizing the Language of a Grammar

S → NP VP
VP → V NP
VP → VP PP
NP → NP PP
PP → P NP
NP → Kim | snow | Oslo
V → saw
P → in

All Complete Derivations

- are rooted in the start symbol S;
- label internal nodes with categories $\in C$, leafs with words $\in \Sigma$;
- instantiate a grammar rule $\in P$ at each local subtree of depth one.



A Simple-Minded Parsing Algorithm

Control Structure

- top-down: given a parsing goal α, use all grammar rules that rewrite α;
- successively instantiate (extend) the right-hand sides of each rule;
- for each β_i in the RHS of each rule, recursively attempt to parse β_i;
- termination: when α is a prefix of the input string, parsing succeeds.

(Intermediate) Results

- Each result records a (partial) tree and remaining input to be parsed;
- complete results consume the full input string and are rooted in S;
- whenever a RHS is fully instantiated, a new tree is built and returned;
- all results at each level are combined and successively accumulated.

The Recursive Descent Parser

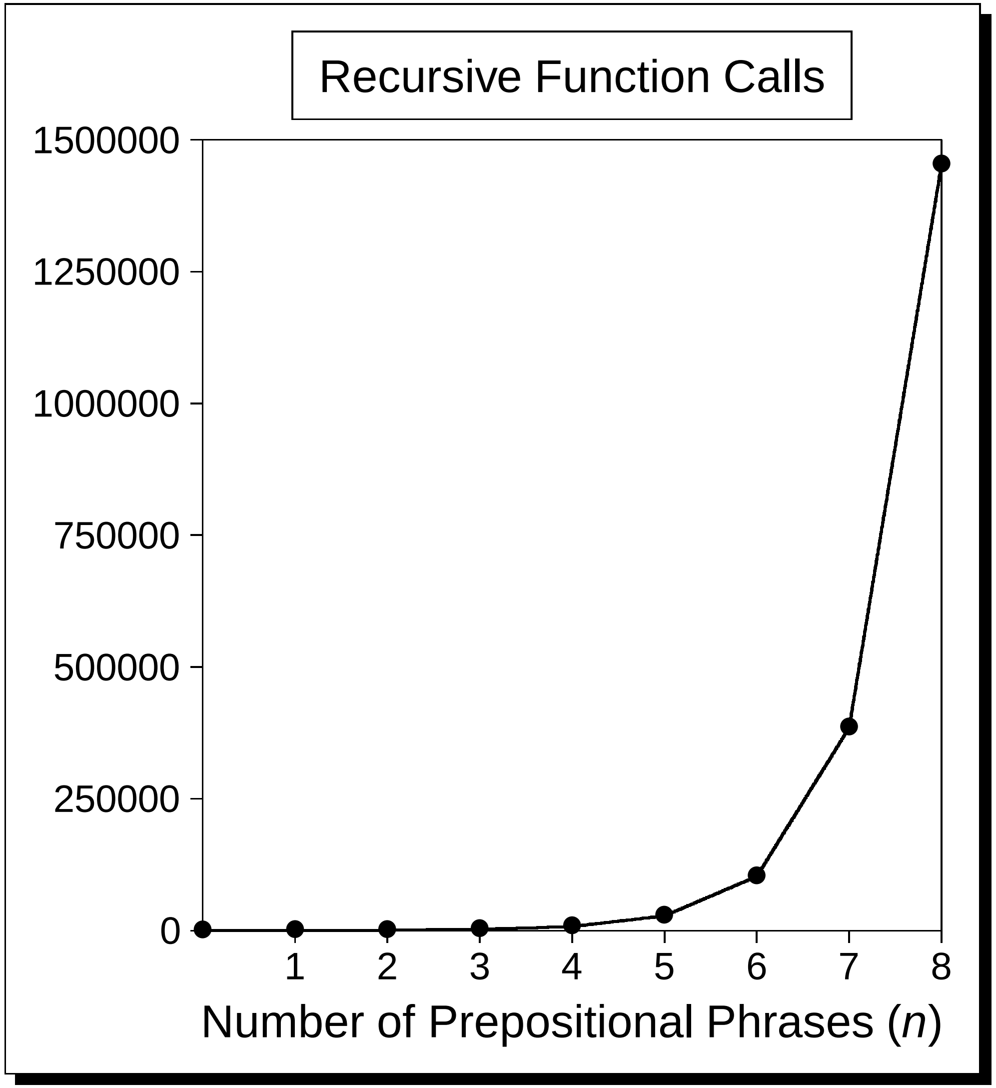
```
(defun parse (input goal)
  (if (equal (first input) goal)
    (let ((edge (make-edge :category (first input))))
      (list (make-parse :edge edge :input (rest input))))
    (loop
        for rule in (rules-deriving goal)
        append (extend-parse (rule-lhs rule) nil (rule-rhs rule) input))))
```

```
(defun extend-parse (goal analyzed unanalyzed input)
  (if (null unanalyzed)
    (let ((edge (make-edge :category goal :daughters analyzed)))
      (list (make-parse :edge edge :input input)))
    (loop
        for parse in (parse input (first unanalyzed))
        append (extend-parse
                goal (append analyzed (list (parse-edge parse)))
                (rest unanalyzed)
                (parse-input parse)))))
```

A Closer Look at the Calling Sequence

```
SSP(18): (parse '(kim adored snow) 's)
parse(): input: (KIM ADORED SNOW); goal: S
  parse(): input: (KIM ADORED SNOW); goal: NP
    parse(): input: (KIM ADORED SNOW); goal: KIM
    parse(): input: (KIM ADORED SNOW); goal: SANDY
    parse(): input: (KIM ADORED SNOW); goal: SNOW
  parse(): input: (ADORED SNOW); goal: VP
    parse(): input: (ADORED SNOW); goal: V
      parse(): input: (ADORED SNOW); goal: LAUGHED
      parse(): input: (ADORED SNOW); goal: ADORED
    parse(): input: (ADORED SNOW); goal: V
      parse(): input: (ADORED SNOW); goal: LAUGHED
      parse(): input: (ADORED SNOW); goal: ADORED
    parse(): input: (SNOW); goal: NP
      ...
```

Quantifying the Complexity of the Parsing Task



Kim adores snow (*in Oslo*)^n

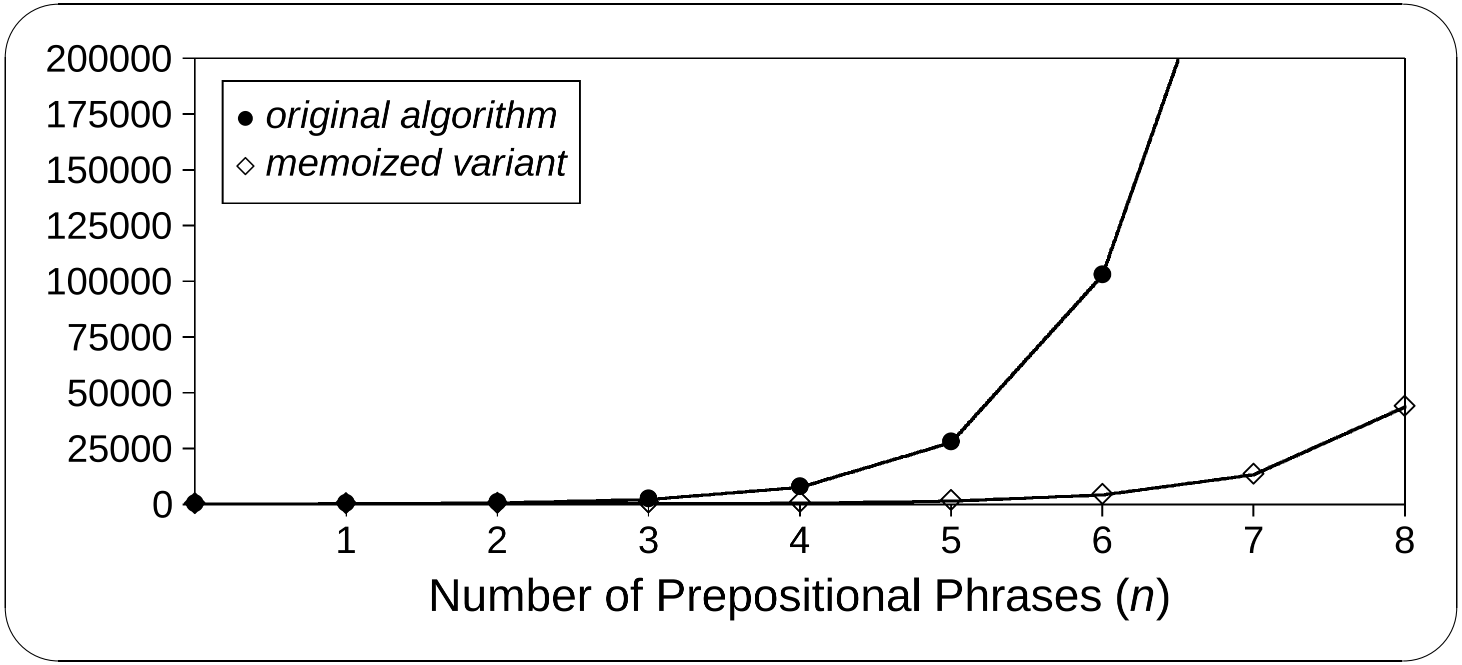
n	**trees**	**calls**
0	1	46
1	2	170
2	5	593
3	14	2,093
4	42	7,539
5	132	27,627
6	429	102,570
7	1430	384,566
8	4862	1,452,776
⋮	⋮	⋮

Memoization: Remember Earlier Results

Dynamic Programming

- The function call `(parse (adored snow) V)` executes two times;
- *memoization* — record `parse()` results for each set of arguments;

$\longrightarrow$ requires abstract data type, efficient indexing on *input* and *goal*.



Top-Down vs. Bottom-Up Parsing

Top-Down (Goal-Oriented)

- Left recursion (e.g. a rule like 'VP $\rightarrow$ VP PP') causes infinite recursion;
- grammar conversion techniques (eliminating left recursion) exist, but will typically be undesirable for natural language processing applications;

$\rightarrow$ assume bottom-up as basic search strategy for remainder of the course.

Bottom-Up (Data-Oriented)

- unary (left-recursive) rules (e.g. 'NP $\rightarrow$ NP') would still be problematic;
- lack of parsing goal: compute all possible derivations for, say, the input *adores snow*; however, it is ultimately rejected since it is not sentential;
- availability of partial analyses desirable for, at least, some applications.

Chart Parsing — Specialized Dynamic Programming

Basic Notions

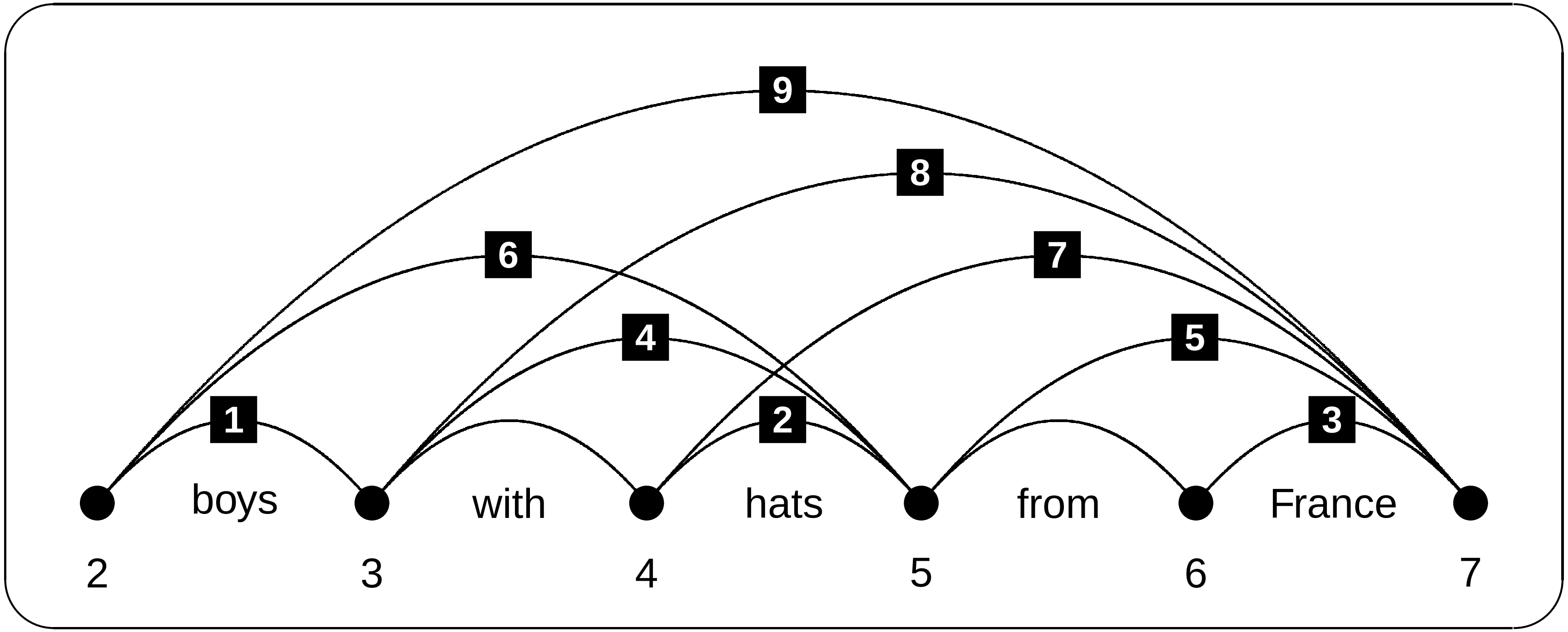
- Use *chart* to record partial analyses, indexing them by string positions;
- count inter-word vertices; CKY: chart row is *start*, column *end* vertex;
- treat multiple ways of deriving the same category for some substring as *equivalent*; pursue only once when combining with other constituents.

Key Benefits

- Dynamic programming (memoization): avoid recomputation of results;
- efficient indexing of constituents: no search by start or end positions;
- compute *parse forest* with exponential 'extension' in *polynomial* time.

Bounding Ambiguity — The Parse Chart

- For many substrings, more than one way of deriving the same category;
- NPs: 1 | 2 | 3 | 6 | 7 | 9; PPs: 4 | 5 | 8; 9 ≡ 1 + 8 | 6 + 5;
- *parse forest* — a single item represents multiple trees [Billot & Lang, 89].



The CKY (Cocke, Kasami, & Younger) Algorithm

for ($0 \leq i < |$*input*$|$) do
  chart$_{[i,i+1]} \leftarrow \{\alpha \,|\, \alpha \rightarrow$ *input*$_i \in P\}$;
for ($1 \leq l < |$*input*$|$) do
  for ($0 \leq i < |$*input*$| - l$) do
    for ($1 \leq j \leq l$) do
      if ($\alpha \rightarrow \beta_1\,\beta_2 \in P \wedge \beta_1 \in$ *chart*$_{[i,i+j]} \wedge \beta_2 \in$ *chart*$_{[i+j,i+l+1]}$) then
        chart$_{[i,i+l+1]} \leftarrow$ *chart*$_{[i,i+l+1]} \cup \{\alpha\}$;

$[0,2] \leftarrow [0,1] + [1,2]$
$\ldots$
$[0,5] \leftarrow [0,1] + [1,5]$
$[0,5] \leftarrow [0,2] + [2,5]$
$[0,5] \leftarrow [0,3] + [3,5]$
$[0,5] \leftarrow [0,4] + [4,5]$

	1	2	3	4	5
0	NP		S		S
1		V	VP		VP
2			NP		NP
3				P	PP
4					NP

Limitations of the CKY Algorithm

Built-In Assumptions

- *Chomsky Normal Form* grammars: $\alpha \rightarrow \beta_1\beta_2$ or $\alpha \rightarrow \gamma$ ($\beta_i \in C$, $\gamma \in \Sigma$);
- breadth-first (aka exhaustive): always compute all values for each cell;
- rigid control structure: bottom-up, left-to-right (one diagonal at a time).

Generalized Chart Parsing

- Liberate order of computation: no assumptions about earlier results;
- *active edges* encode partial rule instantiations, 'waiting' for additional (adjacent and passive) constituents to complete: $[1, 2, \text{VP} \rightarrow \text{V} \bullet \text{NP}]$;
- parser can fill in chart cells in *any* order and guarantee completeness.