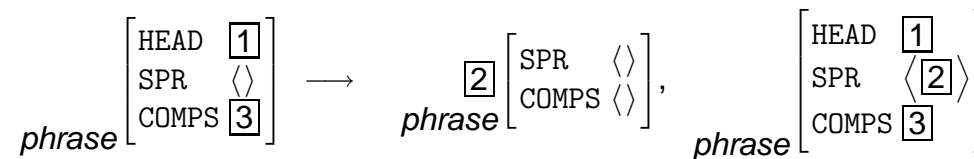


Computational Linguistics (INF2820 — TFSs)



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Limitations of (Our) Context-Free Grammars

Agreement and Valency (For Example)

That dog barks.

**That dogs barks.*

**Those dogs barks.*

The dog chased a cat.

**The dog barked a cat.*

**The dog chased.*

**The dog chased a cat my neighbours.*

The cat was chased by a dog.

**The cat was chased of a dog.*

...



A Really Complicated Language

[...] *omdat ik Jan Henk de nijlpaarden zag helpen voeren .*



More Terminology: Grammatical Functions

Licensing — Government — Agreement

*The dog barks. — *The dog a cat barks — *The dog barks a cat.*

*Kim depends on Sandy — *Kim depends in Sandy*

The class meets on Thursday in 508 at 12:15.

- **Constituent** node in analysis tree (terminal or instantiation of rule);
- **Head** licenses additional constituents and can govern their form;
- **Specifier** precedes head, singleton, nominative case, agreement;
- **Complement** post-head, licensed and governed, order constraints;
- **Adjunct** ‘free’ modifier, optional, may iterate, designated position;
- **Government** directed: a property of c_1 determines the form of c_2 ;
- **Agreement** bi-directional: co-occurrence of properties on c_1 and c_2 .



An Ambiguous Example

Kim shoveled snow on lifts.



A Highly Ambiguous Example

The manager placed his bid on my desk.



Structured Categories in a Unification Grammar

- All (constituent) categories in the grammar are typed feature structures;
 - specific TFS configurations may correspond to ‘traditional’ categories;
- labels like ‘S’ or ‘NP’ are mere abbreviations, not elements of the theory.

word $\left[\begin{array}{l} \text{HEAD } \textit{noun} \\ \text{SPR } \langle \langle \text{HEAD } \textit{det} \rangle \rangle \\ \text{COMPS } \langle \rangle \end{array} \right]$

‘N’

‘lexical’

phrase $\left[\begin{array}{l} \text{HEAD } \textit{verb} \\ \text{SPR } \langle \rangle \\ \text{COMPS } \langle \rangle \end{array} \right]$

‘S’

‘maximal’

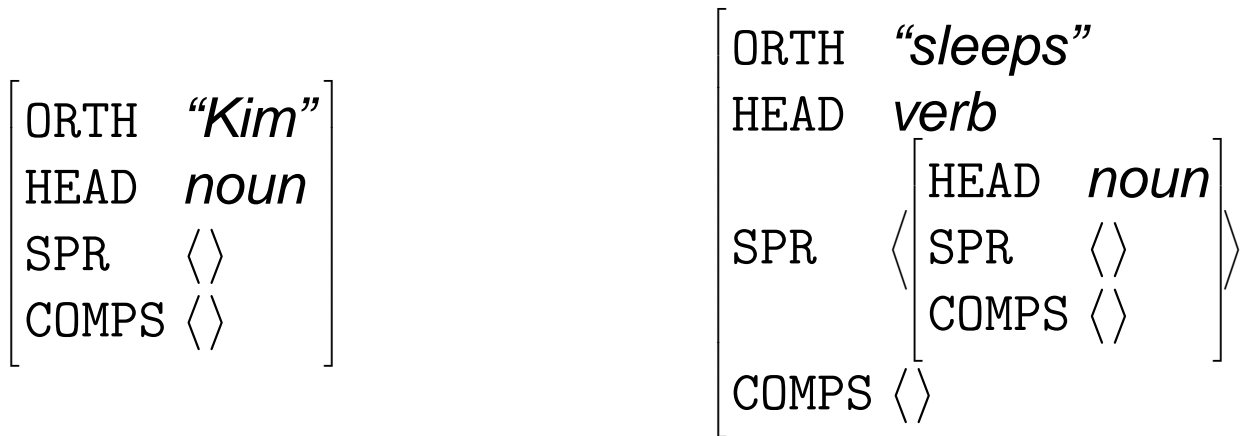
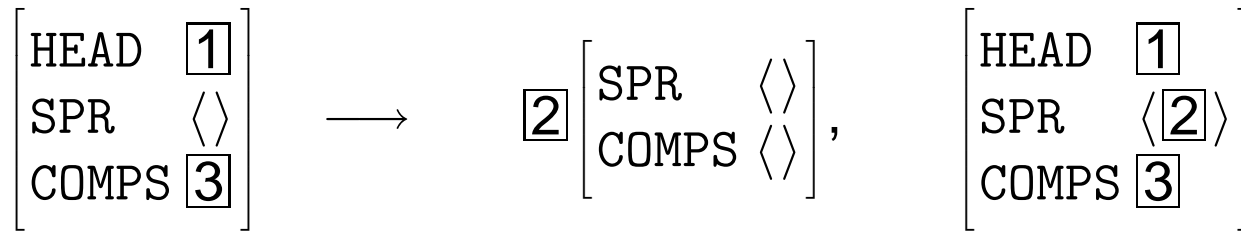
phrase $\left[\begin{array}{l} \text{HEAD } \textit{verb} \\ \text{SPR } \langle \langle \text{HEAD } \textit{noun} \rangle \rangle \\ \text{COMPS } \langle \rangle \end{array} \right]$

‘VP’

‘intermediate’

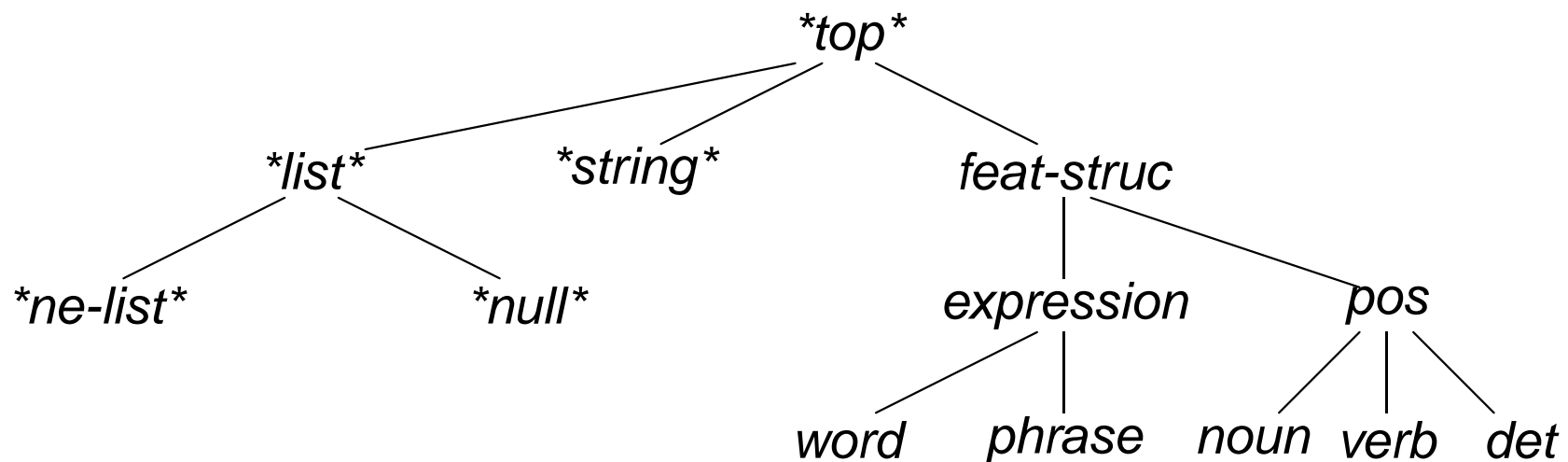


Interaction of Lexicon and Phrase Structure Schemata



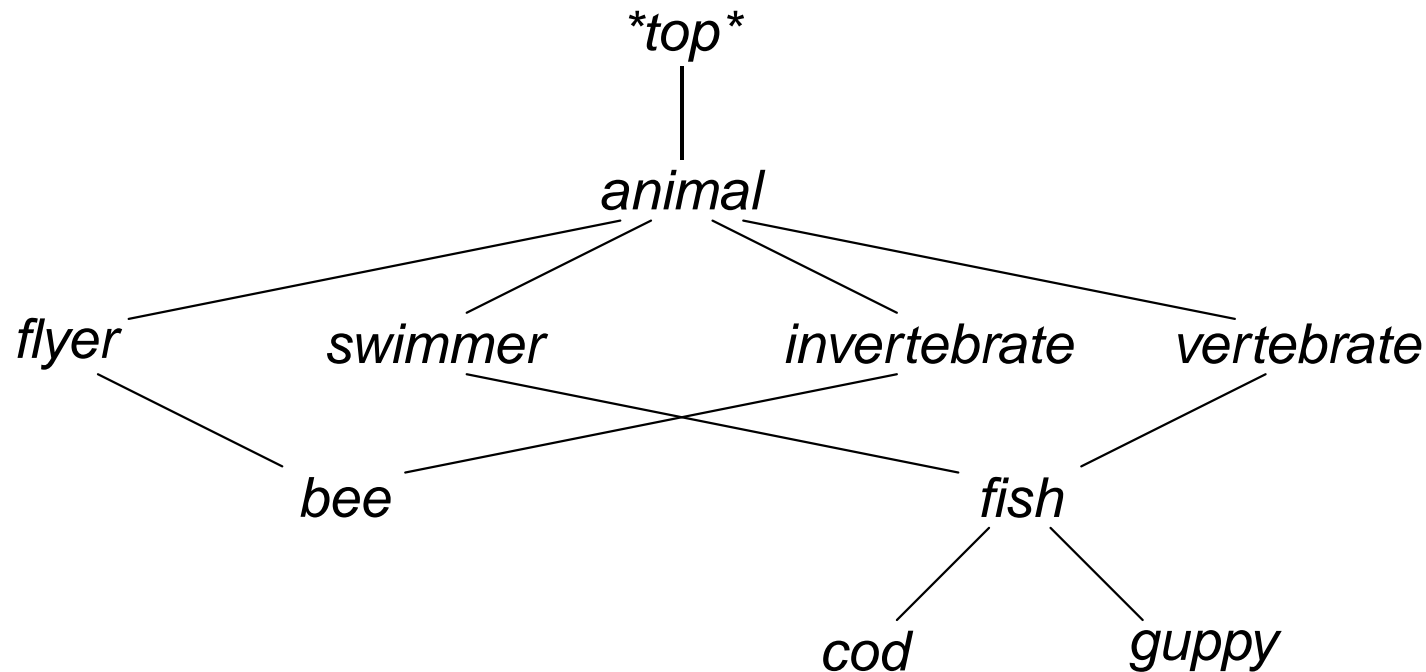
The Type Hierarchy: Fundamentals

- Types ‘represent’ groups of entities with similar properties (‘classes’);
- types ordered by specificity: subtypes inherit properties of (all) parents;
- type hierarchy determines which types are compatible (and which not).



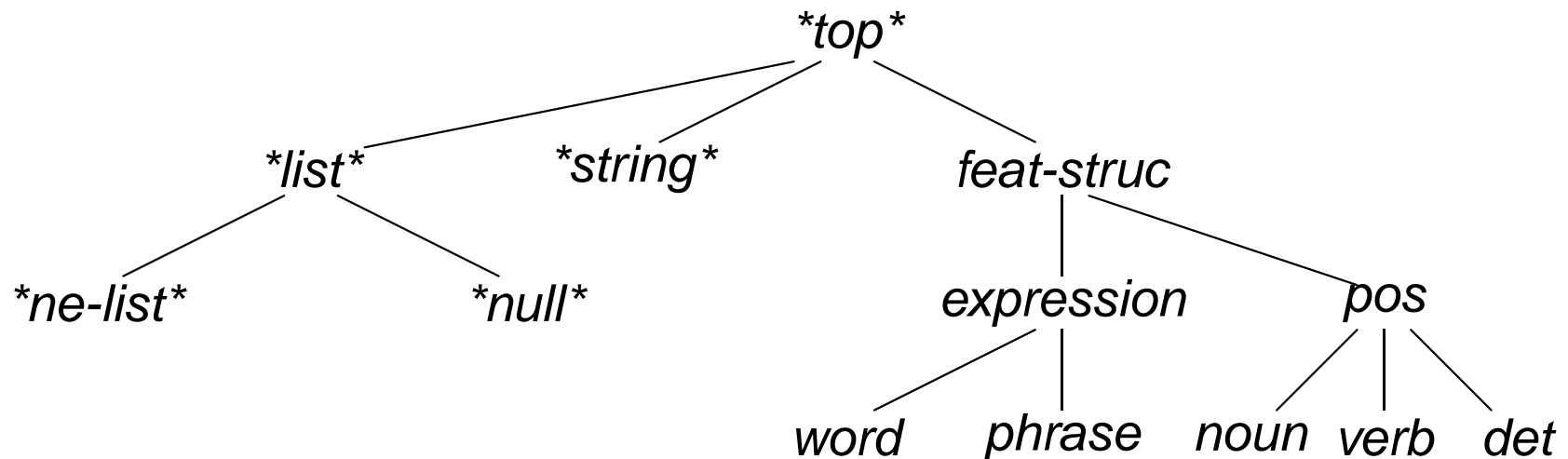
Multiple Inheritance

- *flyer* and *swimmer* no common descendants: they are incompatible;
- *flyer* and *bee* stand in hierarchical relationship: they unify to subtype;
- *flyer* and *invertebrate* have a unique greatest common descendant.



The Type Hierarchy — Appropriate Features

- Features record properties of entities; in turn, feature values are TFSs;
- features are defined by a unique most general type: appropriateness;
- feature values constrained to a specific type → monotonic inheritance.



Typed Feature Structure Subsumption

- Typed feature structures can be partially ordered by information content;
- a more general structure is said to *subsume* a more specific one;
- $*top*$ is the most general feature structure (while \perp is inconsistent);
- \sqsubseteq ('square subset or equal') conventionally used to depict subsumption.

Feature structure F subsumes feature structure G ($F \sqsubseteq G$) iff: (1) if path p is defined in F then p is also defined in G and the type of the value of p in F is a supertype or equal to the type of the value of p in G , and (2) all paths that are reentrant in F are also reentrant in G .



Feature Structure Subsumption: Examples

$$\text{TFS}_1: \begin{matrix} a \\ \left[\begin{array}{l} \text{FOO } x \\ \text{BAR } x \end{array} \right] \end{matrix}$$

$$\text{TFS}_2: \begin{matrix} a \\ \left[\begin{array}{l} \text{FOO } x \\ \text{BAR } y \end{array} \right] \end{matrix}$$

$$\text{TFS}_3: \begin{matrix} b \\ \left[\begin{array}{l} \text{FOO } y \\ \text{BAR } x \\ \text{BAZ } x \end{array} \right] \end{matrix}$$

$$\text{TFS}_4: \begin{matrix} a \\ \left[\begin{array}{l} \text{FOO } \boxed{1} x \\ \text{BAR } \boxed{1} \end{array} \right] \end{matrix}$$

Hierarchy

a	FOO		x
	BAR		
b	BAZ		y

Feature structure F subsumes feature structure G ($F \sqsubseteq G$) iff: (1) if path p is defined in F then p is also defined in G and the type of the value of p in F is a supertype or equal to the type of the value of p in G , and (2) all paths that are reentrant in F are also reentrant in G .



Typed Feature Structure Unification

- Decide whether two typed feature structures are mutually compatible;
- determine combination of two TFSs to give the most general feature structure which retains all information which they individually contain;
- if there is no such feature structure, unification fails (depicted as \perp);
- unification *monotonically* combines information from both 'input' TFSs;
- *relation to subsumption* the unification of two structures F and G is the most general TFS which is subsumed by both F and G (if it exists).
- \sqcap ('square set intersection') conventionally used to depict unification.



Typed Feature Structure Unification: Examples

$$\text{TFS}_1: a \begin{bmatrix} \text{FOO } x \\ \text{BAR } x \end{bmatrix}$$

$$\text{TFS}_2: a \begin{bmatrix} \text{FOO } x \\ \text{BAR } y \end{bmatrix}$$

$$\text{TFS}_3: b \begin{bmatrix} \text{FOO } y \\ \text{BAR } x \\ \text{BAZ } x \end{bmatrix}$$

$$\text{TFS}_4: a \begin{bmatrix} \text{FOO } \boxed{1} x \\ \text{BAR } \boxed{1} \end{bmatrix}$$

Hierarchy

a	FOO		x
	BAR		
b	BAZ		y

$$\text{TFS}_1 \sqcap \text{TFS}_2 \equiv \text{TFS}_2 \quad \text{TFS}_1 \sqcap \text{TFS}_3 \equiv \text{TFS}_3 \quad \text{TFS}_3 \sqcap \text{TFS}_4 \equiv b \begin{bmatrix} \text{FOO } \boxed{1} y \\ \text{BAR } \boxed{1} \\ \text{BAZ } x \end{bmatrix}$$


Type Constraints and Appropriate Features

- Well-formed TFSs satisfy all *type constraints* from the type hierarchy;
- type constraints are typed feature structures associated with a type;
- the top-level features of a type constraint are *appropriate features*;
- type constraints express generalizations over a ‘class’ (set) of objects.

type	constraint	appropriate features
<i>*ne-list*</i>	<i>*ne-list*</i> $\left[\begin{array}{l} \text{FIRST } *top* \\ \text{REST } *list* \end{array} \right]$	FIRST and REST



More Interesting Well-Formed Unification

Type Constraints Associated to *animal* Hierarchy

$$\begin{array}{l}
 \text{swimmer} \rightarrow \text{swimmer} \left[\begin{array}{l} \text{FINS } \textit{bool} \end{array} \right] \quad \text{mammal} \rightarrow \text{mammal} \left[\begin{array}{l} \text{FRIENDLY } \textit{bool} \end{array} \right] \\
 \\
 \text{whale} \rightarrow \text{whale} \left[\begin{array}{l} \text{BALEEN } \textit{bool} \\ \text{FINS } \textit{true} \\ \text{FRIENDLY } \textit{bool} \end{array} \right]
 \end{array}$$

$$\text{mammal} \left[\begin{array}{l} \text{FRIENDLY } \textit{true} \end{array} \right] \sqcap \text{swimmer} \left[\begin{array}{l} \text{FINS } \textit{bool} \end{array} \right] \equiv \text{whale} \left[\begin{array}{l} \text{BALEEN } \textit{bool} \\ \text{FINS } \textit{true} \\ \text{FRIENDLY } \textit{true} \end{array} \right]$$

$$\text{mammal} \left[\begin{array}{l} \text{FRIENDLY } \textit{true} \end{array} \right] \sqcap \text{swimmer} \left[\begin{array}{l} \text{FINS } \textit{false} \end{array} \right] \equiv \perp$$



Recursion in the Type Hierarchy

- Type hierarchy must be finite *after* type inference; illegal type constraint:

```
*list* := *top* & [ FIRST *top*, REST *list* ].
```

- needs additional provision for empty lists; indirect recursion:

```
*list* := *top*.
```

```
*ne-list* := *list* & [ FIRST *top*, REST *list* ].
```

```
*null* := *list*.
```

- recursive types allow for *parameterized list types* ('list of X'):

```
*s-list* := *list*.
```

```
*s-ne-list* := *ne-list* & *s-list* &  
[ FIRST expression, REST *s-list* ].
```

```
*s-null* := *null* & *s-list*.
```



Notational Conventions

- lists not available as built-in data type; abbreviatory notation in TDL:

$$\langle a, b \rangle \equiv [\text{FIRST } a, \text{REST } [\text{FIRST } b, \text{REST } *null*]]$$

- underspecified (variable-length) list:

$$\langle a, \dots \rangle \equiv [\text{FIRST } a, \text{REST } *list*]$$

- difference (open-ended) lists; allow concatenation by unification:

$$\langle ! a ! \rangle \equiv [\text{LIST } [\text{FIRST } a, \text{REST } \#tail], \text{LAST } \#tail]$$

- built-in and ‘non-linguistic’ types pre- and suffixed by asterisk (**top**);
- strings (e.g. “*chased*”) need no declaration; always subtypes of **string**;
- strings cannot have subtypes and are (thus) mutually incompatible.



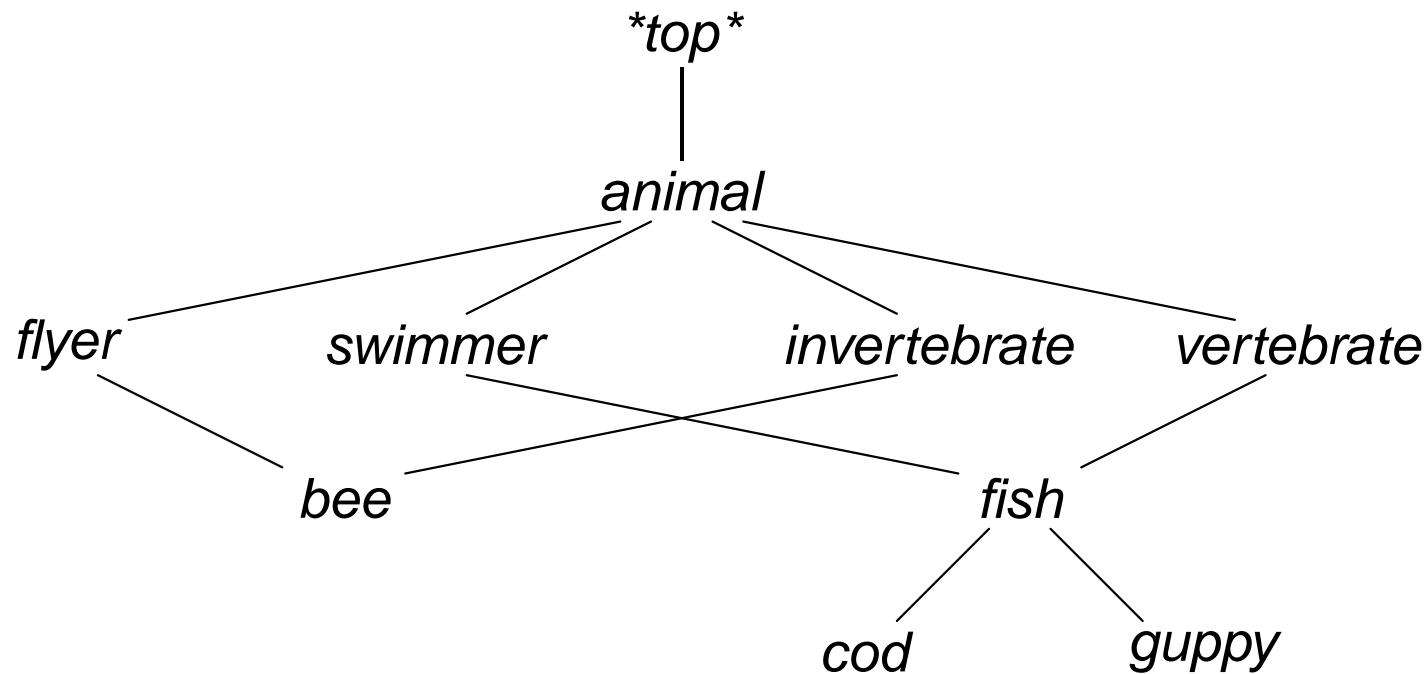
Properties of (Our) Type Hierarchies

- **Unique Top** a single hierarchy of all types with a unique top node;
- **No Cycles** no path through the hierarchy from one type to itself;
- **Unique Greatest Lower Bounds** Any two types in the hierarchy are either (a) incompatible (i.e. share no descendants) or (b) have a unique most general ('highest') descendant (called their greatest lower bound);
- **Closed World** all types that exist have a known position in hierarchy;
- **Compatibility** type compatibility in the hierarchy determines feature structure unifiability: two types unify to their greatest lower bound.



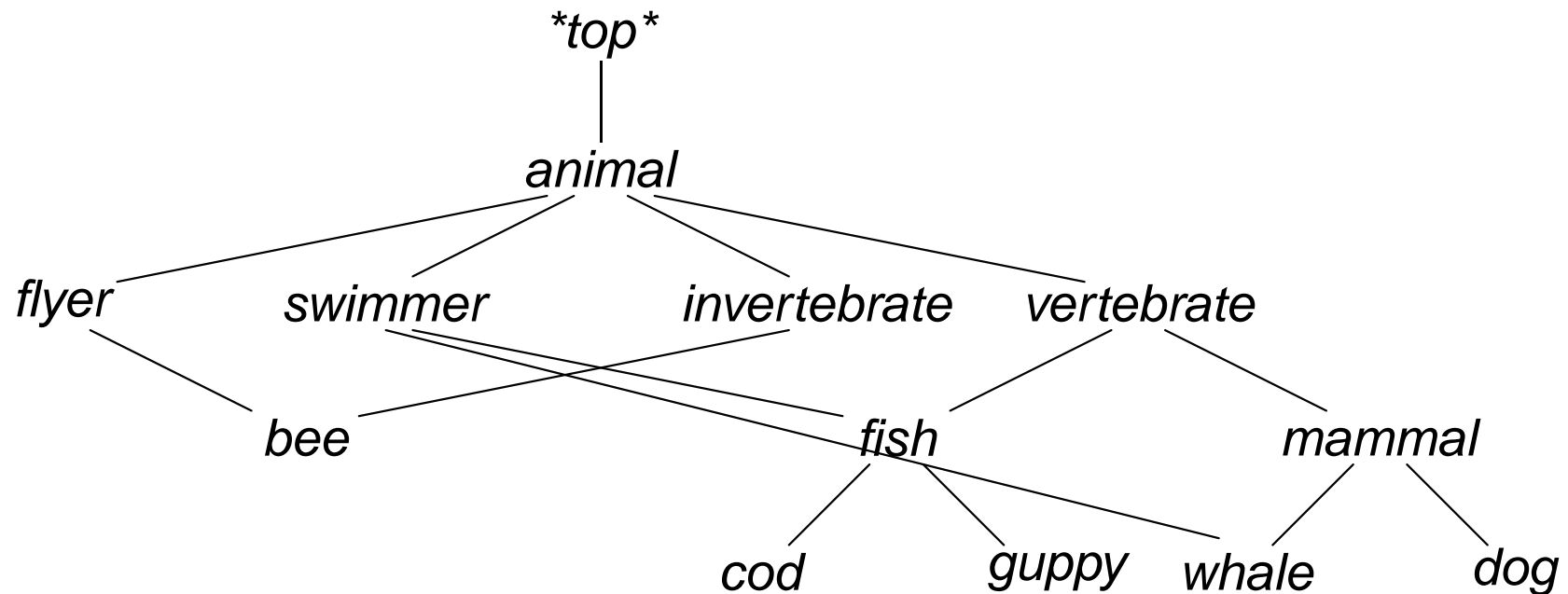
Multiple Inheritance (Repeated for Convenience)

- *flyer* and *swimmer* no common descendants: they are incompatible;
- *flyer* and *bee* stand in hierarchical relationship: they unify to subtype;
- *flyer* and *invertebrate* have a unique greatest common descendant.



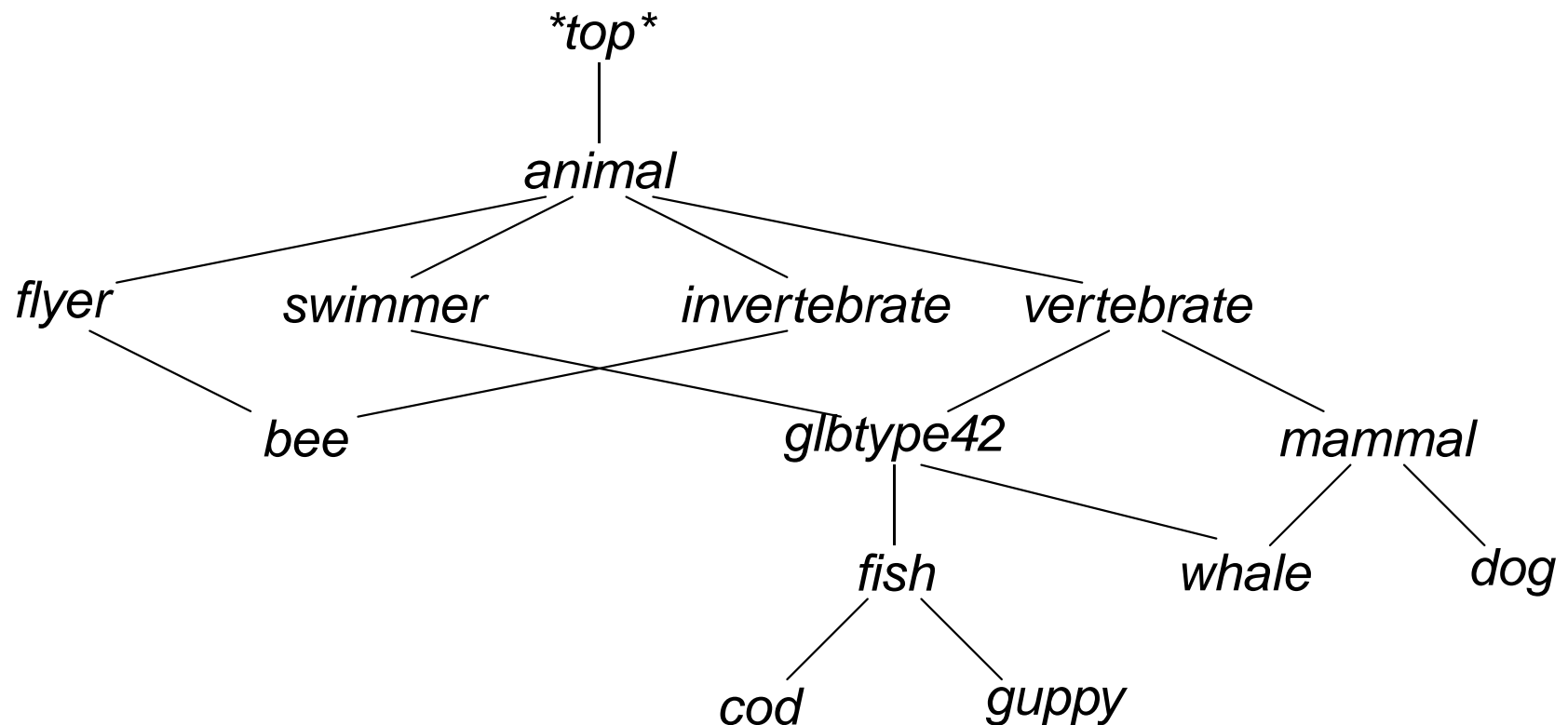
An Invalid Type Hierarchy

- *swimmer* and *vertebrate* have two joint descendants: *fish* and *whale*;
- *fish* and *whale* are incomparable in the hierarchy: glb condition violated.



Fixing the Type Hierarchy

- LKB system introduces *glb* types as required: 'swimmer-vertebrate'.

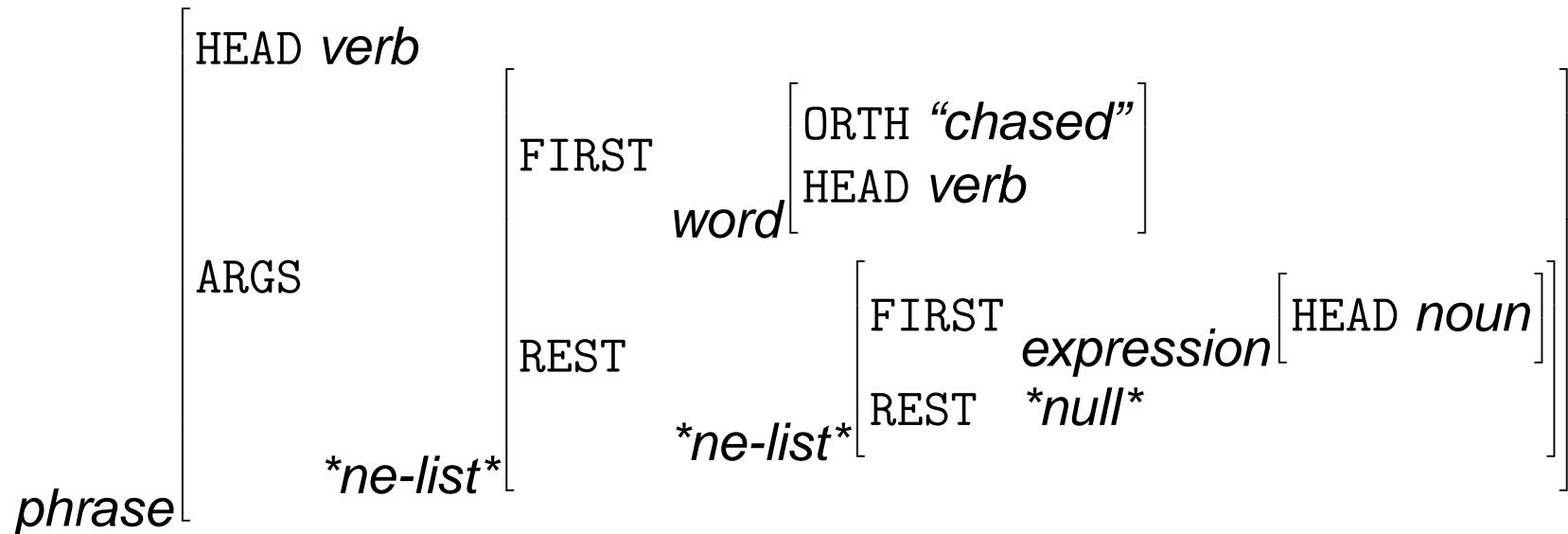


Properties of (Our) Typed Feature Structures

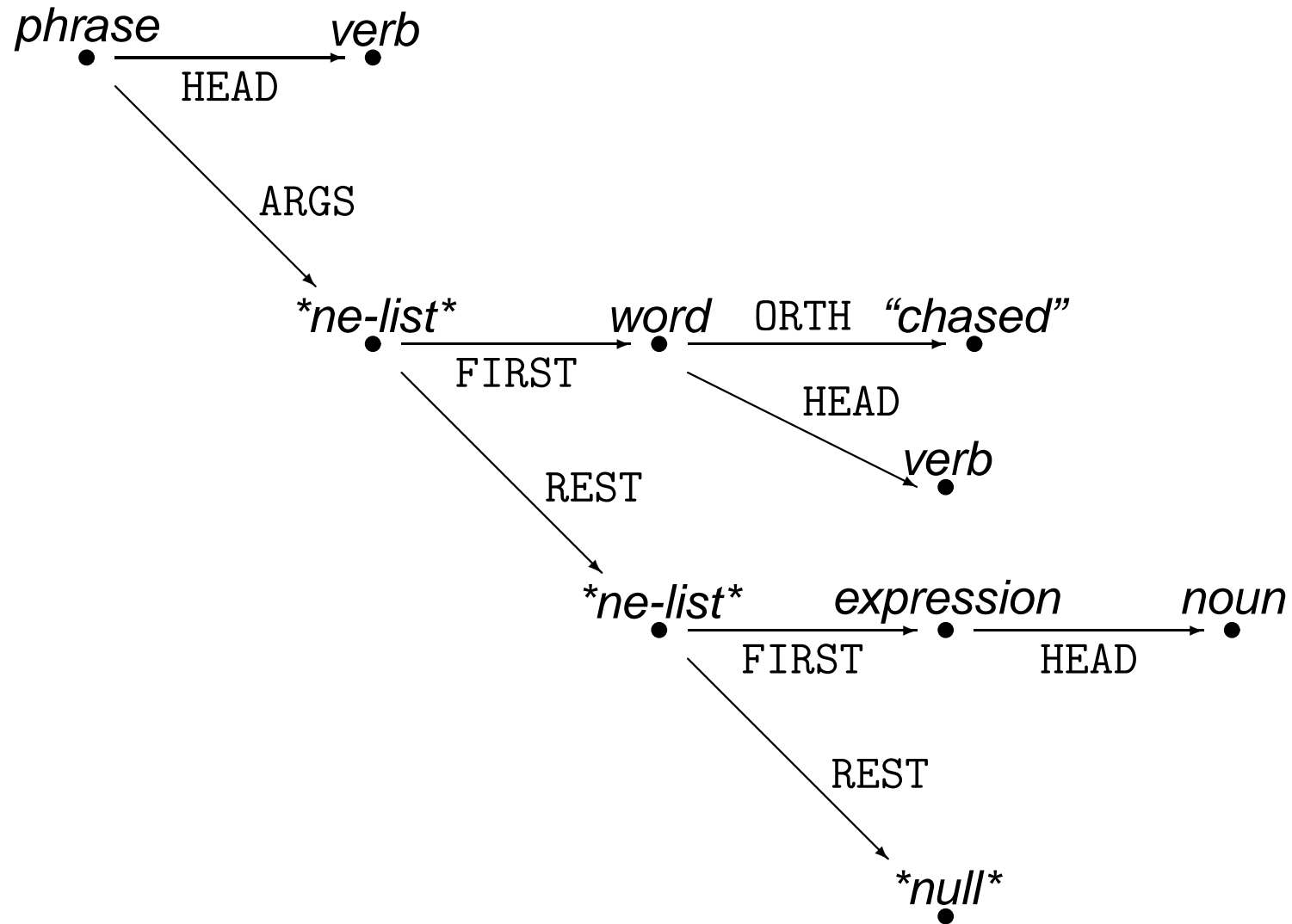
- **Finiteness** a typed feature structure has a finite number of nodes;
- **Unique Root and Connectedness** a typed feature structure has a unique root node; apart from the root, all nodes have at least one parent;
- **No Cycles** no node has an arc that points back to the root node or to another node that intervenes between the node itself and the root;
- **Unique Features** any node can have any (finite) number of outgoing arcs, but the arc labels (i.e. features) must be unique within each node;
- **Typing** each node has single type which is defined in the hierarchy.



Typed Feature Structure Example (as AVM)



Typed Feature Structure Example (as Graph)

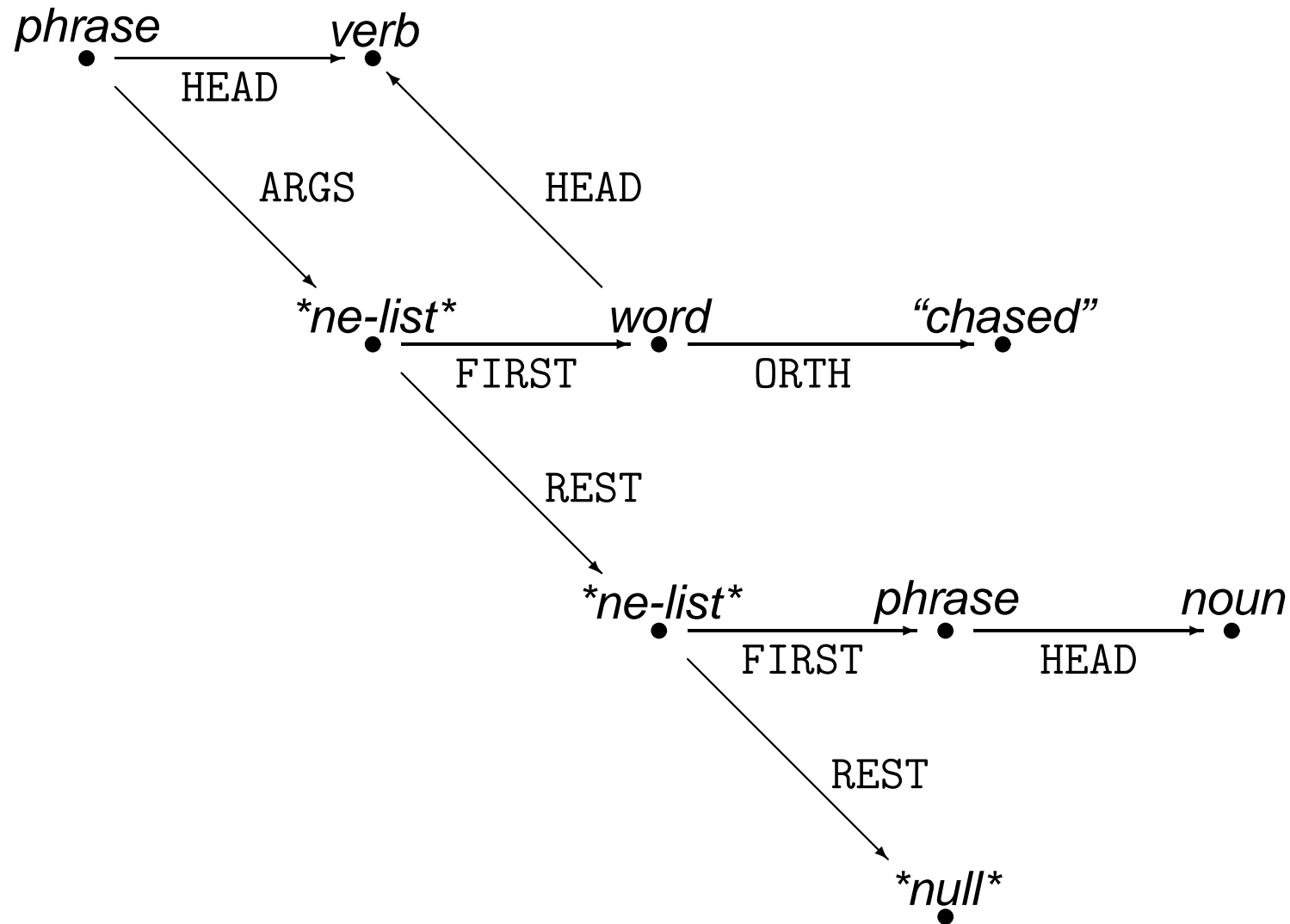


Typed Feature Structure Example (in TDL)

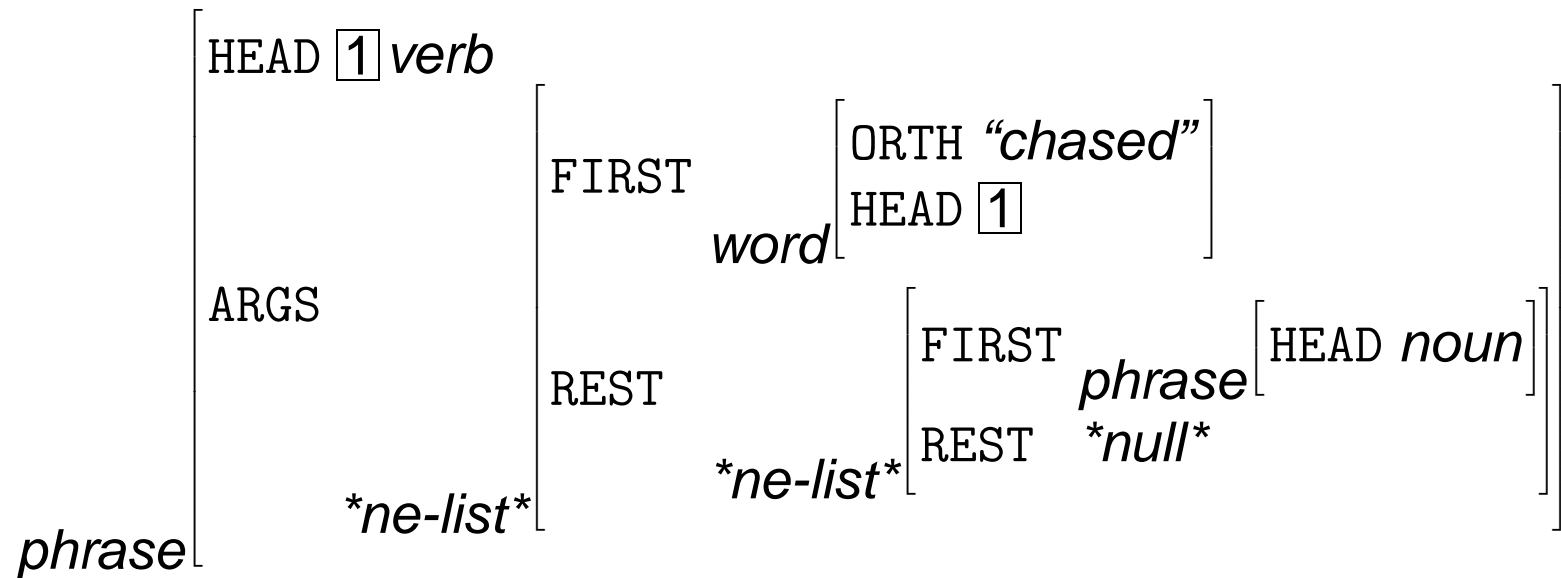
```
vp := phrase &
[ HEAD verb,
  ARGS *ne-list* &
    [ FIRST word &
      [ ORTH "chased",
        HEAD verb ],
      REST *ne-list* &
        [ FIRST expression &
          [ HEAD noun ],
          REST *null* ]]] .
```



Reentrancy in a Typed Feature Structure (Graph)



Reentrancy in a Typed Feature Structure (AVM)



Reentrancy in a Typed Feature Structure (TDL)

```
vp := phrase &
[ HEAD #head & verb,
  ARGS *ne-list* &
    [ FIRST word &
      [ ORTH "chased",
        HEAD #head ],
      REST *ne-list* &
        [ FIRST phrase &
          [ HEAD noun ],
          REST *null* ]]] .
```



The Linguistic Knowledge Builder (LKB)

Compiler and Interactive Debugger

- Grammar definition errors identified at load time by position in file;
- inheritance and appropriateness tracked by type and attributes;
- batch check, expansion, and indexing of full lexicon on demand;
- efficient parser and generator to map between strings and meaning;
- visualization of main data types; interactive stepping and unification.

- Main developers: Copestake (original), Carroll, Malouf, and Oepen;
- implementation: Allegro CL, Macintosh CL, (LispWorks, CMU CL);
- available in open-source and binary form for common platforms.

